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Preprint 1973 No. 2

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January 1973.

1. Introduction and Summary. We shall consider the convex cone \breve{K} of finite continuous convex functions defined on a convex set K in \mathbb{R}^2 . A large class of extremal functions is identified and it is proved that the extremal functions are dense i in \breve{K} .

Thus the results are very different from the results obtained for convex functions on intervals, where all the extremal functions are of the form a v b for some affine functions a and b, see Blaschke and Pick [1].

It is easily seen that

$$\tilde{K} \cap (-\tilde{K}) = \tilde{A},$$

where Á denote the affine functions. Hence the cone is not pointed and the usual definition of an extremepoint has to be modified as follows:

<u>Definition</u>: Let f, g, and h be elements of K, then f is called extremal if for all g and h such that

$$f = \frac{1}{2}(g+h)$$

there exist a constant $\alpha > 0$ and an affine function a, such that

$$f = \alpha g + a$$
.

We shall apply the following concepts from the theory of convex sets and functions, see Rockafellar [2].

A polyhedral set is a closed convex set which is the intersection of a finite number of halfspaces.

A polytope is a compact polyhedral set.

A face of a polyhedral set P is a subset $F \subset P$ with the property that

$$x \in F$$
, $y \in P$, $z \in P$, $x = \frac{1}{2}(y+z) \Rightarrow y \in F$, $z \in F$.

The O dimensional faces are the extreme points or the vertices. The 1-dimensional faces are the edges. A polyhedral set is bounded by a finite number of edges and has a finite number of vertices. A polytope is the convex hull of its vertices.

The family of polyhedral sets P_1, \ldots, P_m is called a covering of the convex set K if

$$K \subset P_1 \cup \cdots \cup P_m$$

and

$$riP_i \cap riP_i = \emptyset, i \neq j.$$

Here ri denotes the relative interior. The order of a vertex of a polyhedral set in the covering is the number of polyhedral sets which contains it.

Let a_1 , ... a_n denote a family of affine functions, then

is called a polyhedral function and it is seen that it is convex and continuous, and that the sets

$$P_{i} = \{f=a_{i}\}, dim P_{i} = 2,$$

give rise to a covering of R^2 , and therefore of K, by polyhe-dral sets.

The class of extremal functions on R^2 which is found here can be described as polyhedral functions which induce a covering where the vertices are of order 3, see Theorem 1.

An alternative way of describing them is as follows: consider the cylinder

$$C = \{(x, \mu)\}: \mu \ge 0, x \in \mathbb{R}^2\}.$$

Any affine function $a \in A$ can be thought of as cutting away from C the set

$$\{(x,\mu) : 0 \le \mu \le a(x), x \in \mathbb{R}^{2}\}$$

After having cut C by means of the affine functions a_1, \dots, a_n we are left with the epigraph of $f = 0 \vee a_1 \vee \dots \vee a_n$

$$C_n = \{(x,\mu) : \mu \ge f(x), x \in \mathbb{R}^2\}.$$

If each function a_m , m = 1, 2, ..., n is chosen such that it does not cut through an extremepoint of C_{m-1} then f will be extreme, since then the vertices of the covering induced by f will have order 3.

The result that the functions thus constructed can approximate any continuous convex function f uniformly on a compact set is now rather obvious since the epigraph of f can be cut out of C by continuing the above procedure, each time avoing the vertices already created, see Theorem 2.

The actual proofs for convex sets in R² are more complicated since we need extra conditions to ensure that there are enough vertices inside K, see the lemma.

2. A Combinatorial Lemma.

Lemma. Let C be an open convex set of dimension 2. Let P_1 , ..., P_m be a covering of C with convex closed polyhedral sets of dimension 2 such that

- 1) Each polyhedral set has a vertex in C.
- 2) Any two vertices in C can be connected by a path of edges in C.
- 3) No vertex is in the relative interior of an edge.
- 4) Each vertex is of order 3.

Let finally f be a continuous function which is affine on each P_i , i = 1,...,m. Then if f = 0 on two polyhedral sets. with a common edge then f = 0 on C.

<u>Proof</u>: Let P_1 and P_2 have a common edge $L = P_1 \cap P_2$ and let f = 0 on $P_1 \cup P_2$. We want to prove that there is a third polyhedral set P_3 which has an edge L_1 in common with P_1 and an

edge L_2 in common with P_2 , then since f is affine on P_3 and 0 on L_1 and L_2 , the condition 3) will ensure that f = 0 on P_3 .

If L \cap C contains no vertex of P₁ (or P₂) then L would bisect C C, such that P₁ and P₂ would be on different sides of L. Now P₁ and P₂ each have a vertex in C₀ by condition 1), and by 2) they can be connected by a path of edges inside C. This path must meet L \cap C and hence L \cap C does contain a vertex. Let therefore V₁ be a vertex in L \cap C. Since V₁ is of order 3 there is a third polyhedral set P₃ which meets P₁ and P₂ at V₁. This set clearly has a common edge with P₁ and with P₂ and by the above argument f = 0 on P₃.

This was the start of the induction and we can now prove that the conditions 1) through 4) are sufficient that the presented argument spreads to all of C.

Let us assume that f = 0 on

 $C_{k} = P_{1} \cup \dots \cup P_{k}, (k \le m),$

which is connected and contains a vertex in the interior, namely V_1 .

If $C \sim C_k = \emptyset$ we have proved that f = 0 on C. Otherwise let $x \in C \sim C_k$, and let P be a polyhedral set which contains x. Let V_2 be a vertex in C \cap P.

Now consider the connected path that leads from V_1 to V_2 in π^- side C. Let V_3 be the first vertex on this path which lies on the boundary of C_k .

Notice that since V_3 is reach from inside C_k , the three edges that meet at V_3 all lie in C_k , but since V_3 is on the boundary of C_k and in C there must be a polyhedral set P_{k+1} say, which is outside C_k and which meets the boundary of C_k at V_3 and at two edges. Hence since f is affine on P_{k+1} and 0 on C_k we find that f is 0 on P_{k+1} as well. This completes the induction and the proof of the lemma. Let us remark, that in case $C = R^2$ there are two cases possible. Either there are no vertices at all in which case the covering (P_1, \ldots, P_m) consists of parallel strips covering the plane, or there is at least one vertex, in which case 1) and 2) are automatically satisfied.

Notice also that if the covering is induced by a convex function then 3) is automatically satisfied and hence we see that the most important condition is the fourth condition that each vertex should have order 3.

3. The Extremal Functions. We can not identify all the extremal convex functions, but we can find so many that there a are enough to prove the main result that they are dense in \check{K} .

<u>Theorem 1.</u> A polyhedral convex function f is extremal in \tilde{K} if the covering of K given by f satisfies the conditions 1), 2), and 4) for C = int K.

Proof: Let g and h be elements of K and let

$$f = \frac{1}{2}(g + h)$$
.

Let P_1 , ..., P_m denote the covering induced by f which satisfies the conditions 1),2) and 4). The condition 3) will then automatically be satisfied.

Clearly g and h must be affine on each of the polyhedral sets P_1 , ..., P_m and we shall assume that g and h are polyhedral functions.

Now let a, b, and c denote affine functions, such that

$$f_1 = f-a$$
, $g_1 = g-b$, and $h_1 = h-c$

all vanish on P1, then

$$f_1 = \frac{1}{2}(g_1 + h_1).$$

Let then $x_0 \in int P_2 \cap int K$, where P_2 has an edge in common with P_1 .

If $g_1(x_0) = 0$ then g_1 is 0 on P_2 and P_1 and by the lemma $g_1 \equiv 0$ on K which proves that g is affine and hence that f is extremal.

Let us therefore assume that $g_1(x_0) > 0$ and by a similar argument that $h_1(x_0) > 0$.

Let us then define

$$f_2 = f_1/f_1(x_0), g_2 = g_1/g_1(x_0), h_2 = h_1/h_1(x_0)$$

then

$$f_2 = \alpha g_2 + (1 - \alpha) h_2$$

where

$$0 < \alpha = g_1(x_0)/2f_1(x_0) < 1.$$

Now we have that

$$f_2 = g_2 = h_2 = 0 \text{ on } P_1$$

and

$$f_2(x_0) = f_2(x_0) = h_2(x_0) = 1$$

but then

$$f_2 = g_2 = h_2 \text{ on } P_2.$$

If we apply the lemma to the piecewise affine function $f_2^{-g_2}$ we get that

 $f_2 = g_2 = h_2 \text{ on } K$

which implies that f is extremal. This completes the proof of Theorem 1.

<u>Proposition</u>: Let a and b be affine functions then a and aVb are extremal functions. If c is an affine function such that the equations

$$a(x) = b(x) = c(x)$$

have only one solution in int K then a v b v c is extremal.

<u>Proof</u>: It is easily seen that a is extremal, and that the construction in the proof of Theorem 1 will give that $a \lor b$ is extremal. The above condition on c ensures the existence of a vertex in int K and the covering induced by $a \lor b \lor c$ satisfies the conditions 1), 2), and 4) of the lemma.

<u>Corollary</u>. In the convex cone of finite continuous convex functions defined on R^2 , the polyhedral functions which induce coverings with vertices of order 3 are extremal.

In particular the functions a, $a \vee b$, and $a \vee b \vee c$ are extremal if the equations

$$a(x) = b(x) = c(x)$$

have only one solution.

<u>Proof</u>: This follows from the remarks after the proof of the lemma, together with Theorem 1 and the proposition.

We shall now prove the main result.

<u>Theorem 2</u>. Any finite continuous convex function on the con + evex set K can be approximated uniformly on any convex compact subset of K by an extremal convex function.

<u>Proof</u>: Let $f \in K$ be given as a finite continuous convex function on K. Let $K_1 \subset K$ be a compact convex set of dimension 2.

We shall prove that f can be approximated by modifying the function on K_1 a finite number of times in such a way that the final modification is an extremal function in \tilde{K} , which differs 1 less than ϵ from f on K_1 .

1) The first modification is to approximate f by a polyhedral function

$$\begin{array}{c} f \\ 1 \\ 1 \\ 1 \\ 1 \\ i \\ n \end{array}$$

as follows: For each $x \in K_1$ we find a subgradient a_x and determine a neighbourhood N_x of x, such that $f(y) < a_x(y) + \epsilon/4$, $y \in N_x$.

By compactness we can pick out a finite number of neighbourhoods which cover K_1 and the corresponding subgradients provide us with the function f_1 .

2) The next step consists in modifying f_1 such that the polyhedral covering induced by f_1 satisfies condition 1) of the lemma. Let therefore P be such a polyhedral set where f_1 is affine and such that

int P
$$\cap$$
 int K₁ $\neq \phi$.

The set P need not have any vertices in int K_1 but let us choose

 $x_1 \in int \mathbb{P}P \cap int K_1$.

Now choose three affine functions a, b, and c, such that

$$a(x) = b(x) = c(x)$$

only have the solution $x = x_1$. The function

f₁ + δ a v b v c

will be a convex polyhedral function with a vertex at x_1 . This function clearly induces a covering of K_1 with more polyhedral sets than before, but each new polyhedral set will have a vertex in int K_1 .

We then repeat this construction for each polyhedral set from f_1 which does not posses a vertex in int K_1 . The final modification f_2 will consist of f_1 plus a sum of simple extremal convex functions and will have an induced covering satisfying the first condition of the lemma. For δ sufficiently small $|f_2-f_1| < \varepsilon/4$ on K_1 .

3) The vertices of f_2 need not be connected but let x_1 and x_2 be any two vertices in int K_1 . Let a be an affine func-

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tion such that

$$a(x_1) = a(x_2) = 0$$

and such that $a \neq 0$.

The function

 $f_2 + \delta(0va)$

is a convex polyhedral function whose covering of K₁ will contain some new polyhedral sets. Each new set however, will have a vertex on the line determined by x₁ and x₂ inside int K₁, and x₁ and x₂ can be connected by a path of edges in int K₁, also lying on the line [x₁, x₂]. This procedure is repeated by replacing x₂ by any of the vertices from f₂ each time adding an a simple extremal function. We end up with a function f₃ whose covering satisfies conditions 1) and 2) of the lemma. For δ sufficiently small we get $|f_2 - f_3| < \epsilon/4$ on K₁.

4) The final step consists in ensuring that all vertices have order 3.

Assume V to be a vertex of order s > 3 and let a be a subgradient such that

 $f_{3}(x) > a(x), x \neq V$ $f_{3}(x) = a(x), x = V$

Then

 $f_3 v (a+\delta)$

is a convex polyhedral function with the property that for δ sufficiently small the corresponding polyhedral covering will be changed only around V in such a way that V will be surrounded by a small polytope with s vertices each of order 3.

Clearly the new polyhedral sets constructed this way still have vertices in int K_1 and these can still be connected in-

side int K₁.

This construction is repeated for each vertex of order > 3 each time taking the maximum of the function so far obtained and a suitable affine function. The final function will induce a covering with all the desired properties listed in the lemma.

For δ sufficiently small this function f_4 will lie within $\varepsilon/4$ of f_3 on K_1 .

Thus we can apply Theorem 1 and we get that f_4 is an extremal function which differs less than \in from f on K_1 as was to be proved.

4. References.

- [1] Blaschke, W. and Pick, G., Distanzschätzungen im Funktionenraum II. Math. Ann. 77, 277-302 (1916).
- [2] Rockafellar, R.T., Convex Analysis, Princeton University Press (1970).