





From Truth, Belief and Knowledge to Tsallis Entropy

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indications: key notions, key questions

truth
belief, tendency to act
knowledge, perception, extended experience
interaction
experiment, preparation, control
description, effort, information

what is “information”? – a potential saving of effort!
what then is “entropy”? – minimal effort, given the truth!
and, what *can* we know? – well, depends on your belief ...



information: cost and associated effort

What is the **cost of information** or, how much are you willing to pay – or *have* to pay – in order to know that an event has happened?

Or, what is the **effort** you are willing to/have to allocate?

Depends on the probability t , you *believe* the event has: $\kappa(t)$.
 κ tells us the **individual effort**. It is the **effort-function** or the **descriptor**.

effort \longleftrightarrow description ?

Requirements: $\kappa(1) = 0$, κ is smooth (and decreasing).

Further, natural with **normalization** via the **differential cost**
 $\iota = -\kappa'(1)$. If $\iota = 1$, we obtain **natural units**, nats;
 if $\iota = \ln 2$, we measure in **binary units**, bits.



the power hierarchy, the exponential hierarchy

Which kind of descriptors would you expect?

Note that any “reasonable” monotone function f defines a descriptor via linearization. Simply take

$$\kappa(t) = \frac{f(1) - f(t)}{f'(1)}.$$

Suggestions: The **power hierarchy** is defined from the functions $t \mapsto t^a$ ($-\infty < a < \infty$) and gives the descriptors:

$$t \mapsto \frac{1 - t^a}{a}.$$

... and the **exponential hierarchy** is defined from the functions $t \mapsto b^t$ ($b > 0$) and gives the descriptors

$$t \mapsto \frac{1 - b^{t-1}}{\ln b}.$$

But are any of these “sensible”? – and what does that mean?



accumulated effort in a probabilistic context

Consider distributions over a discrete **alphabet** \mathbb{A} :

$x = (x_i)_{i \in \mathbb{A}}$ represents **truth**, $y = (y_i)_{i \in \mathbb{A}}$ represents **belief**.

Accumulated effort (expected per observation) is

$$\Phi(x, y) = \sum_{i \in \mathbb{A}} x_i \kappa(y_i).$$

Φ satisfies the **perfect match principle** (for short is **proper**) if

$$\Phi(x, y) \geq \Phi(x, x)$$

with equality only for $y = x$ (or $\Phi(x, x) = \infty$).

Theorem There is only one descriptor κ , the **classical descriptor**, for which Φ above is proper, viz. (nats)

$$\kappa(t) = \ln \frac{1}{t}.$$



questioning the basic definition

Surely, $\Phi(x, y) = \sum x_i \kappa(y_i)$ is the right expression for accumulated effort **as seen by someone, who knows the truth as well as the belief** ...

... but is this how **you perceive** accumulated effort?

What if the x_i 's above are not what you perceive as truth?

... perhaps this also depends on what you believe – and Φ should rather be something like $\sum \pi(x_i, y_i) \kappa(y_i)$.

Let's leave the probabilistic setting for a while and go philosophical:



the beginnings of a philosophy of information

The whole is the **world**, \mathcal{V}

Situations from the world involve **Nature** and you, **Observer**.

Nature has no **mind** but holds the **truth** (x),

Observer has a **creative** mind,

- seeks the **truth** (x)
- is confined to **belief** (y)
- aims at **knowledge** (z).

But what is “knowledge”? Knowledge is

- the **synthesis of extensive experience**
- an expression of how Observer **perceives** situations from \mathcal{V}
- a **manifestation** of truth for Observer, for you.



interaction and effort

Proposal: Knowledge (z) depends on truth and belief via a characteristic **interactor** Π : $z = \Pi(x, y)$ $\mathcal{V} = \mathcal{V}_\Pi$.

$\Pi_1 : (x, y) \mapsto x$ defines the **classical world** \mathcal{V}_1

$\Pi_0 : (x, y) \mapsto y$ defines a **black hole** \mathcal{V}_0

$\Pi_q : (x, y) \mapsto qx + (1 - q)y$ defines mixtures, **Tsallis'** \mathcal{V}_q 's

Associated with \mathcal{V}_Π are, possibly many, effort functions, Φ 's. Extending the previous definition, an effort function is **proper** if it satisfies the **perfect match principle** (PMP):

$\Phi(x, y) \geq \Phi(x, x)$ with equality iff there is a **perfect match**, i.e. $y = x$ (or $\Phi(x, x) = \infty$).

Thesis Given \mathcal{V}_Π , there is at most one proper Φ -function



entropy, divergence, the fundamental inequality

Abstract modelling involves **effort** (Φ), **entropy** (H), and **divergence** (D). Φ is assumed proper.

Entropy is defined as **minimal effort, given the truth**, divergence as **excess (or redundant) effort**:

$$H(x) = \Phi(x, x); \quad D(x, y) = \Phi(x, y) - H(x).$$

The properness of Φ may be expressed in terms of D by **the fundamental inequality of information theory (FI)**:

$$D(x, y) \geq 0 \quad \text{with equality iff } y = x.$$

Further notions and properties are best discussed for probabilistic modelling.



probabilistic modelling (discrete)

Truth-, belief- and knowledge instances are $x = (x_i)$, $y = (y_i)$ and $z = (z_i)$ (i ranging over an alphabet \mathbb{A}).

x and y are probability distributions, z just a function on \mathbb{A} .

Interaction, Π , acts via the **local interactor** π :

$(\Pi(x, y))_{i \in \mathbb{A}} = (\pi(x_i, y_i))_{i \in \mathbb{A}}$. π is always assumed **sound**, i.e. $\pi(s, t) = s$ if $t = s$ (perfect match).

π is **weakly consistent** if $\forall x \forall y : \sum z_i = 1$. **Strong consistency** requires that z is always a probability distribution.

Proposition: Only the π_q 's given by $\pi_q(s, t) = qs + (1 - q)t$ are weakly consistent; strong consistency requires $0 \leq q \leq 1$.



accumulated effort, the one and only

Accumulated effort always chosen among $\Phi_{\pi, \kappa}$ where κ is a descriptor and

$$\Phi_{\pi, \kappa}(x, y) = \sum_{i \in \mathbb{A}} \pi(x_i, y_i) \kappa(y_i).$$

Theorem (modulo regularity conditions). Given $\pi = \pi(s, t)$, let $\pi'_2 = \frac{\partial \pi}{\partial t}$ and put $\chi(t) = \pi'_2(t, t)$.

Only one among the $\Phi_{\pi, \kappa}$'s can be proper, viz. the solution to

$$t\kappa'(t) + \chi(t)\kappa(t) = -1, \quad \kappa(1) = 0. \quad (*)$$

If π is consistent, hence one of the π_q 's, then a proper $\Phi_{\pi, \kappa}$ exists iff $q > 0$ ($q = 0$ OK as a singular case, though).

If so, the unique descriptor concerned is the one depending linearly on t^{q-1} , i.e. $\kappa_q(t) = \ln_q \frac{1}{t}$ (recall: $\ln_q u = \frac{1}{1-q}(u^{1-q} - 1)$).



gross effort, pointwise fundamental inequality

Introduce **gross (accumulated) effort** and **gross entropy** by adding a term representing **overhead cost** (or effort):

$$\text{gross effort: } \tilde{\Phi}(x, y) = \sum_{i \in \mathbb{A}} (\pi(x_i, y_i) \kappa(y_i) + y_i) = \Phi(x, y) + 1,$$

$$\text{gross entropy: } \tilde{H}(x) = \sum_{i \in \mathbb{A}} (x_i \kappa(x_i) + x_i) = H(x) + 1.$$

Clearly, “gross divergence” = divergence and, defining the **divergence generator** by

$\delta(s, t) = (\pi(s, t) \kappa(t) + t) - (s \kappa(s) + s)$, one has $D(x, y) = \sum \delta(x_i, y_i)$. We refer to the inequality $\delta \geq 0$ as the **pointwise fundamental inequality** (PFI). Clearly

PFI \implies FI. **Conjecture** **Converse also true**

In practice, PMP and FI are always proved via PFI !



given κ , which world are you in?

Given π , we insist, when possible, to choose κ such that the resulting function Φ is proper. This gives a unique choice, the **ideal descriptor**. You determine κ from π , but

Warning: you cannot determine π from κ

Thus knowing the entropy function does not reveal the world.

Examples: Let $\pi = \pi_q$ ($q > 0$) and consider π^ξ of the form

$$\pi^\xi(s, t) = \xi^{-1} \left(\pi(\xi(s), \xi(t)) \right).$$

Then the differential equation (*) is unchanged, hence you find the same descriptor κ_q . E.g. for $\xi(u) = \ln u$, $\pi^\xi(s, t) = s^q t^{1-q}$; by PFI, the associated effort is proper.

Problem which κ 's are associated with (meaningful) π 's?

e.g. $\kappa(t) = \frac{1}{2}(t^{-2} - 1)$; or $\kappa(t) = 1 - \exp(t - 1)$?



what can we know? (abstract modelling)

Setting: World \mathcal{V}_π with ideal descriptor and effort fct. Φ .

I.J. Good (1952): **Belief is a tendency to act !**

To us, this is expressed via **controls**, w 's. There is a bijection $y \leftrightarrow w$ ($w = \hat{y}$; $y = \check{w}$). In our probabilistic modelling this is given by $w_i = \kappa(y_i)$; $i \in \mathbb{A}$.

Expressed via controls, the effort function is denoted Ψ :
 $\Psi(x, w) = \Phi(x, y)$ with $y \leftrightarrow w$.

What can Observer do? via control ! preparations which are sets of x 's, typically denoted by \mathcal{P} .

A **feasible preparation** is one which Observer can **realize**.



more on preparations

Typical example (of **genus 1**): Fix a control w and a **level** h .
Set-up an experiment (!?) which constrains Nature's possibilities to the preparation

$$\mathcal{P}(w, h) = \{x \mid \Psi(x, w) = h\}$$

or to the variant $\mathcal{P}_{\leq}(w, h) = \{x \mid \Psi(x, w) \leq h\}$.

Finite non-empty intersections of such **level sets**
(or **sub-level sets**) constitute the feasible preparations and
show what Observer can know !



equilibrium, MaxEnt and all that!

A closer study of a fixed preparation \mathcal{P} requires **game theory** and exploits thinking of John Nash. We shall only outline this. The two players are Nature with truth instances $x \in \mathcal{P}$ as strategies and Observer with controls w as strategies. As **objective function** we take $\Psi = \Psi(x, w)$. The maxmin value is easily seen to be the **MaxEnt value**

$$H_{\max}(\mathcal{P}) = \sup_{x \in \mathcal{P}} H(x).$$

If this equals the minmax value (required finite), the game is in **equilibrium**.

Another notion, often overlooked: A control ε^* is **robust** if, for some finite h , $\Psi(x, \varepsilon^*) = h$ for all $x \in \mathcal{P}$. Then h is the **level of robustness**.

By results of Nash:



robustness lemma, exponential families

Robustness lemma If $x^* \in \mathcal{P}$ and $\varepsilon^* = \hat{x}^*$ is robust with level h , then the game is in equilibrium with x^* and ε^* as **optimal strategies**, in particular, x^* is the MaxEnt strategy. (furthermore, the celebrated **Pythagorean inequalities** hold).

Let w be a control, let \mathcal{L}^w be the **preparation family** of non-empty sets of the form $\mathcal{P}(w, h)$. The associated **exponential family**, denoted $\hat{\mathcal{E}}^w$ is the set of controls ε which are robust for all preparations in \mathcal{L}^w . From robustness lemma:

Consider a preparation family \mathcal{L}^w . Let x^* be a truth instance, put $\varepsilon^* = \hat{x}^*$ and assume that $\varepsilon^* \in \hat{\mathcal{E}}^w$. Put $h = \Psi(x^*, w)$. Then the game corresponding to $\mathcal{P}(w, h)$ is in equilibrium and has x^* and ε^* as optimal strategies. In particular, x^* is the MaxEnt distribution for $\mathcal{P}(w, h)$.



sketch of MaxEnt calculations in \mathcal{V}_q

Return to probabilistic setting and consider a Tsallis world $\mathcal{V} = \mathcal{V}_q$, cor. to π_q with $q > 0$.

Fix $y \longleftrightarrow w$. Then \mathcal{L}^w consists of all preparations \mathcal{P} for which $\Psi(x, w)$ is constant over \mathcal{P} .

But $\Psi(x, w) = \sum (qx_i + (1 - q)y_i)w_i$ so condition is equivalent to $\sum x_i w_i$ being constant over \mathcal{P} .

For fixed constants α and β this implies that $\sum x_i(\alpha + \beta w_i)$ is constant over \mathcal{P} .

Now, if $\alpha + \beta w$ is a control, say w^* , $\sum x_i w_i^*$ is constant over \mathcal{P} , hence $\Psi(x, w^*)$ is constant over \mathcal{P} , i.e. $w^* \in \hat{\mathcal{E}}^w$ and the robustness lemma applies.

Then, given β , try to adjust α so that $\alpha + \beta w$ is a control.

Classically, α is the logarithm of the **partition function**.

Finally, adjust β (\approx inverse temperature) to desired level ...



what have we achieved?

- found a reasonably transparent interpretation of Tsallis entropy
- developed a basis for an abstract theory
- clarified role of FI via PMP; focus on PFI as the natural basis for establishing FI and hence PMP
- identified the unit of entropy as an overhead
- answered the question “what *can* we know”
- found good (*the right* ?) definition of an exponential family
- indicated dual role of preparations and exponential families
- exploited games and wisdom of Nash, enabled MaxEnt calculations without introducing Lagrange multipliers
- separated Nature from Observer in key expressions



what needs being done?

- interaction, how?
- description, how?
- control, how?
- expand, quantum setting ...
- link to information geometry
- ...

thank you !

