Homotopy colimits in model categories

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1 Introduction

In [1], Dwyer and Spalinski construct the so-called homotopy pushout functor, motivated by the following observation. In the category **Top** of topological spaces, one can construct the *n*-dimensional sphere S^n by glueing together two *n*-disks D^n along their boundaries S^{n-1} , i.e. by the pushout of

$$D^n \stackrel{i}{\longleftrightarrow} S^{n-1} \stackrel{i}{\longrightarrow} D^n$$

where i denotes the inclusion. Let * be the one point space. Observe, that one has a commutative diagram

where all vertical maps are homotopy equivalences, but the pushout of the bottom row is the one-point space * and therefore not homotopy equivalent to S^n . One probably prefers the homotopy type of S^n . Having this idea of calculating the prefered homotopy type in mind, they equip the functor category $\mathbf{C}^{\mathbf{D}}$, where \mathbf{C} is a model category and $\mathbf{D} = a \leftarrow b \rightarrow c$ the category consisting out of three objects a, b, c and two non-identity morphisms as indicated, with a suitable model category structure. This enables them to construct out of the pushout functor colim : $\mathbf{C}^{\mathbf{D}} \rightarrow \mathbf{C}$ its so-called total left derived functor \mathbf{L} colim : $\mathrm{Ho}(\mathbf{C}^{\mathbf{D}}) \rightarrow \mathrm{Ho}(\mathbf{C})$ between the corresponding homotopy categories, which defines the homotopy pushout functor.

Dwyer and Spalinski further indicate how to generalize this construction to define the so-called homotopy colimit functor as the total left derived functor for the colimit functor colim : $\mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ in the case, where \mathbf{D} is a so-called very small category. The goal of this paper is to give a proof of this generalization, since in [1] it is omitted to the reader. The main work lies in proving our Theorem 2, which equips $\mathbf{C}^{\mathbf{D}}$ with the suitable model category structure. For the existence of the total left derived functor for colim, we will use a result from [1].

The paper contains four sections. In section 2 and 3 we recall some definitions and results related to colimits and model categories, respectively. The introduced terminology will be used in section 4, where we construct the homotopy colimit functor.

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2 Colimits

In this section let \mathbf{C} be a category, let \mathbf{D} be a small category and $F : \mathbf{D} \to \mathbf{C}$ a functor. Mainly to fix notations, we recall some definitions and results related to colimits.

Definition 2.1. The functor category $\mathbf{C}^{\mathbf{D}}$, also called the category of diagrams in \mathbf{C} with the shape of \mathbf{D} , is the category, where the objects are functors $\mathbf{D} \to \mathbf{C}$ and the morphisms are natural transformations.

Example 2.2. If **D** is the category $a \leftarrow b \rightarrow c$ with three objects a, b, c and two non-identity morphisms $a \leftarrow b, b \rightarrow c$, then an object X of $\mathbf{C}^{a \leftarrow b \rightarrow c}$ is just a diagram $X(a) \leftarrow X(b) \rightarrow X(c)$ in **C** and a morphism from X to Y in $\mathbf{C}^{a \leftarrow b \rightarrow c}$ is given by a triple (s_a, s_b, s_c) of morphisms in **C** making

$$\begin{array}{c|c} X(a) & \longleftarrow & X(b) \longrightarrow X(c) \\ & \downarrow^{s_a} & \downarrow^{s_b} & \downarrow^{s_c} \\ Y(a) & \longleftarrow & Y(b) \longrightarrow Y(c) \end{array}$$

commute.

Example 2.3. If **D** is the category $a \to b$ with two objects a, b and one nonidentity morphism $a \to b$, then an object of $\mathbf{C}^{a\to b}$ is a morphism $f: X(a) \to X(b)$ in **C** and a morphism from $f: X(a) \to X(b)$ to $g: Y(a) \to Y(b)$ in $\mathbf{C}^{a\to b}$ is just a pair of morphisms (s_a, s_b) in **C** making

$$\begin{array}{ccc} X(a) & \stackrel{f}{\longrightarrow} X(b) \\ & \downarrow^{s_a} & \downarrow^{s_b} \\ Y(a) & \stackrel{g}{\longrightarrow} Y(b) \end{array}$$

commute. We call $\mathbf{C}^{a\to b}$ the *category of morphisms in* C and denote it by $\mathbf{Mor}(\mathbf{C})$.

Example 2.4. If **D** is the category **1** consisting only out of one object 1 and one morphism, then $\mathbf{C}^{\mathbf{D}}$ is isomorphic to **C** via the functor given by $X \mapsto X(1)$ on objects and $f \mapsto f_1$ on morphisms.

Definition 2.5. The constant diagram functor $\Delta = \Delta_{\mathbf{D}} : \mathbf{C} \to \mathbf{C}^{\mathbf{D}}$ is the functor, which sends an object C of **C** to the functor $\Delta(C)$ given by

 $d \mapsto C$ on objects and $g \mapsto id_C$ on morphisms,

and which sends a morphism f of \mathbf{C} to the natural transformation $\Delta(f)$ given by $\Delta(f)_d = f$ in an object d of \mathbf{D} .

Note that a functor $j : \mathbf{D}' \to \mathbf{D}$ from a small category \mathbf{D}' to \mathbf{D} induces a functor $(\cdot)|_{\mathbf{D}',j} = (\cdot)|_{\mathbf{D}'} = (\cdot)|_j : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}^{\mathbf{D}'}$, which sends an object X of $\mathbf{C}^{\mathbf{D}}$ to $X \circ j$ and a morphism $f : X \to Y$ in $\mathbf{C}^{\mathbf{D}}$ to the natural transformation $f|_{\mathbf{D}'} : X|_{\mathbf{D}'} \to Y|_{\mathbf{D}'}$ given by $f_{j(d')}$ in an object d' of \mathbf{D}' . Furthermore, if $j' : \mathbf{D}'' \to \mathbf{D}'$ is a functor from a small category \mathbf{D}'' to \mathbf{D}' , then

$$(\cdot)|_{j \circ j'} = (\cdot)|_{j'} \circ (\cdot)|_j.$$

$$(2.1)$$

Example 2.6. The functor $\mathbf{D} \to \mathbf{1}$ induces the functor $(\cdot)|_{\mathbf{1}} : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}^{\mathbf{D}}$, which composed with the isomorphism $\mathbf{C} \cong \mathbf{C}^{\mathbf{1}}$ yields $\Delta_{\mathbf{D}}$. By (2.1), it follows that $\Delta_{\mathbf{D}'} = (\cdot)|_j \Delta_{\mathbf{D}}$ for any functor j from a small category \mathbf{D}' to \mathbf{D} .

Definition 2.7. A colimit C = (C, t) for $F : \mathbf{D} \to \mathbf{C}$ is an object C of \mathbf{C} together with a natural transformation $t : F \to \Delta(C)$ such that for any object X of \mathbf{C} and any natural transformation $s : F \to \Delta(X)$ there exists a unique morphism $s' : C \to X$ in \mathbf{C} satisfying $\Delta(s')t = s$.

Example 2.8. If $\mathbf{D} = a \leftarrow b \rightarrow c$, then a colimit for an object X of $\mathbf{C}^{a \leftarrow b \rightarrow c}$ is just a pushout of the diagram $X(a) \leftarrow X(b) \rightarrow X(c)$.

Remark 2.9. If a colimit for F exists, we will sometimes speak of the colimit for F for the following reason. If (C, t) and (C', t') are colimits for F, then the unique morphism $h: C \to C'$ in \mathbb{C} such that $\Delta(h)t = t'$ holds, is an isomorphism, which will be called the canonical isomorphism. Given furthermore an object X of \mathbb{C} and a natural transformation $s: F \to \Delta(X)$, let $s'': C' \to X$ be the unique morphism in \mathbb{C} satisfying $\Delta(s'')t' = s$. Then the unique morphism $s': C \to X$ satisfying $\Delta(s')t = s$ is given by s''h.

Remark 2.10. Assume that for any functor $G : \mathbf{D} \to \mathbf{C}$ the colimit $(\operatorname{colim}(G), t_G)$ exists. Then the chosen colimits yield a functor

colim :
$$\mathbf{C}^{\mathbf{D}} \to \mathbf{C}$$
,

called *colimit functor*, which maps a morphism $s : G \to G'$ in $\mathbf{C}^{\mathbf{D}}$ to the unique morphism $s' : \operatorname{colim}(G) \to \operatorname{colim}(G')$ in \mathbf{C} such that $\Delta(s')t_G = t_{G'}s$. Furthermore, we have an adjunction (colim, Δ, α), where the natural equivalence α is given in an object (G, X) of $(\mathbf{C}^{\mathbf{D}})^{\operatorname{op}} \times \mathbf{C}$ by the bijection

$$\alpha = \alpha_{G,X} : \operatorname{Hom}_{\mathbf{C}}(\operatorname{colim}(G), X) \to \operatorname{Hom}_{\mathbf{C}^{\mathbf{D}}}(G, \Delta(X)),$$
$$s' \mapsto \Delta(s')t_G.$$
(2.2)

In particular, any two colimit functors $\mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ are naturally isomorphic and therefore, we will sometimes speak of the colimit functor. **Remark 2.11.** Let \mathbf{D}' be a small category such that there exists an isomorphism $J: \mathbf{C}^{\mathbf{D}} \to \mathbf{C}^{\mathbf{D}'}$ with $J \circ \Delta_{\mathbf{D}} = \Delta_{\mathbf{D}'}$. Assume there exists a colimit functor colim : $\mathbf{C}^{\mathbf{D}'} \to \mathbf{C}$. Then the composition colim $\circ J: \mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ is a colimit functor. Furthermore, for any object X of \mathbf{C} and any natural transformation $s: F \to \Delta_{\mathbf{D}}(X)$ the induced morphism from the colimit of F to X is given by the morphism $\operatorname{colim}(J(F)) \to X$ induced by $J(s): J(F) \to \Delta_{\mathbf{D}'}(X)$.

In the definition of a model category we will use both, the notion of colimit and limit.

Definition 2.12. A limit L = (L, t) for $F : \mathbf{D} \to \mathbf{C}$ is an object L of \mathbf{C} together with a natural transformation $t : \Delta(L) \to F$ such that for any object X of \mathbf{C} and any natural transformation $s : \Delta(X) \to F$ there exists a unique morphism $s' : X \to L$ in \mathbf{C} satisfying $t\Delta(s') = s$.

The following result is proved on page 115 in [3].

Proposition 2.13. Let \mathbf{D}' be a small category and $G : \mathbf{D}' \to \mathbf{C}^{\mathbf{D}}$ a functor. Denote by $\operatorname{ev}_d : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ the evaluation functor in the object d of \mathbf{D} . Assume that for all objects d of \mathbf{D} a limit for the composition $\operatorname{ev}_d \circ G : \mathbf{D}' \to \mathbf{C}$ exists. Then there exists a limit for G.

The notion of colimit is dual to the notion of limit:

Remark 2.14. A colimit for F is the same as a limit for the dual functor $F^{\text{op}} : \mathbf{D}^{\text{op}} \to \mathbf{C}^{\text{op}}$. More precisely, sending a colimit (C, t) for F to (C, t'), where the natural transformation $t' : \Delta_{\mathbf{D}^{\text{op}}}(C) \to F^{\text{op}}$ is given by $(t_d)^{\text{op}}$ in an object d of \mathbf{D}^{op} , defines a one-to-one correspondence between colimits for F and limits for F^{op} .

We conclude this section with a result about pushouts.

Proposition 2.15. Assume that

$$\begin{array}{cccc} X & \stackrel{i}{\longrightarrow} W & and & Y & \stackrel{i'}{\longrightarrow} P \\ & & & & & \\ y & & & & & \\ Y & \stackrel{i'}{\longrightarrow} P & & & Z & \stackrel{i''}{\longrightarrow} Q \end{array}$$

are pushout squares in C. Then so is

$$\begin{array}{ccc} X & \stackrel{i}{\longrightarrow} W \\ & & & & \\ & & & \\ k_j & & & \\ Z & \stackrel{i''}{\longrightarrow} Q \end{array} \end{array} .$$

Proof. From the commutativity of the first two squares one concludes that k'j'i = i''kj, thus the third one commutes too. Given now any commutative square



one has to show, that there is a unique morphism $t : Q \to V$ such that tk'j' = r and ti'' = s. Using that P is a pushout, let $t' : P \to V$ be the unique morphism such that t'j' = r and t'i' = sk. Using that also Q is a pushout, define $t : Q \to V$ to be the unique morphism such that tk' = t' and ti'' = s. This gives the desired morphism t. To show the uniqueness, let $t'' : Q \to V$ be a morphism with t''k'j' = r and t''i'' = s. By the universal property of P, one concludes t''k' = t'. Hence, by the universal property of Q follows t'' = t.

3 Model categories and homotopy categories

3.1 Model categories

The following terms will be used in the definition of a model category.

Definition 3.1. A morphism $f: X \to X'$ of a category **C** is called a *retract* of a morphism $g: Y \to Y'$ of **C**, if there exists a commutative diagram

$$\begin{array}{c} X \xrightarrow{i} Y \xrightarrow{r} X \\ \downarrow f & \downarrow g & \downarrow f \\ X' \xrightarrow{i'} Y' \xrightarrow{r'} X' \end{array}$$

such that $ri = id_X$ and $r'i' = id_{X'}$.

Definition 3.2. Let **C** be a category.

i) Given a commutative diagram in \mathbf{C} of the form

$$\begin{array}{ccc} A \xrightarrow{f} X \\ \downarrow i & \downarrow p \\ B \xrightarrow{g} Y \end{array} \tag{3.1}$$

a *lift* in the diagram is a morphism $h: B \to X$ such that hi = f and ph = g.

ii) Let $i : A \to B$, $p : X \to Y$ be morphisms in **C**. We say that *i* has the *left lifting property* (LLP) with respect to *p* and that *p* has the *right lifting property* (RLP) with respect to *i*, if there exists a lift in any commutative diagram of the form (3.1).

Now we are ready for the definition of a model category.

Definition 3.3. A model category is a category \mathbf{C} together with three classes of morphisms of \mathbf{C} , the class of weak equivalences $W = W(\mathbf{C})$, of fibrations $Fib = Fib(\mathbf{C})$ and of cofibrations $Cof = Cof(\mathbf{C})$, each of which is closed under composition and contains all identity morphisms of \mathbf{C} , such that the following five conditions hold:

MC1: Every functor from a finite category to C has a limit and a colimit.

MC2: If f and g are morphisms of **C** such that gf is defined and if two out of the three morphisms f, g and gf are weak equivalences, then so is the third.

MC3: If f is a retract of a morphism g of **C** and g is a weak equivalence, a fibration or a cofibration, then so is f.

MC4:

- i) Every cofibration has the LLP with respect to all $p \in W \cap Fib$.
- ii) Every fibration has the RLP with respect to all $i \in W \cap Cof$.

MC5: Any morphism f of \mathbf{C} can be factored as

- i) f = pi, where $i \in Cof$, $p \in W \cap Fib$, and as
- ii) f = pi, where $i \in Cof \cap W$, $p \in Fib$.

By a model category structure for a category \mathbf{C} we mean a choice of three classes of morphisms of \mathbf{C} such that \mathbf{C} together with these classes is a model category. Let \mathbf{C} be a model category until the end of this subsection. The following two remarks follow immediately from the definition of a model category.

Remark 3.4. Since any isomorphism $f : X \to Y$ of **C** is a retract of id_Y and since $id_Y \in W \cap Fib \cap Cof$, it follows by **MC3** that also $f \in W \cap Fib \cap Cof$.

Remark 3.5. The opposite category \mathbf{C}^{op} is a model category by defining a morphism f^{op} in \mathbf{C}^{op} to be in $W(\mathbf{C}^{\text{op}})$ if f is in $W(\mathbf{C})$, to be in $Fib(\mathbf{C}^{\text{op}})$ if f is in $Cof(\mathbf{C})$ and to be in $Cof(\mathbf{C}^{\text{op}})$ if f is in $Fib(\mathbf{C})$.

The proofs of the following two propositions can be found on page 87 and 88 of [1].

Proposition 3.6. A morphism i of C is

- i) in Cof, if and only if it has the LLP with respect to all $p \in W \cap Fib$,
- ii) in $W \cap Cof$, if and only if it has the LLP with respect to all $p \in Fib$.

- i) If i is in Cof(C), then so is i'.
- ii) If i is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$, then so is i'.

The next two results are concerned with cofibrations and pushouts and will be used in the proof of Theorem 2.

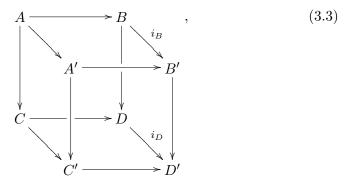
Lemma 3.8. Given a commutative square in C of the form

$$\begin{array}{ccc} A \longrightarrow B \\ \downarrow & \downarrow \\ C \longrightarrow D \end{array} \tag{3.2}$$

Let P, Q be pushouts of $C \leftarrow A \rightarrow B$ and let $i_P : P \rightarrow D$, $i_Q : Q \rightarrow D$ be the morphisms induced by (3.2). Then i_P is a cofibration [resp. weak equivalence] if and only if i_Q is a cofibration [resp. weak equivalence].

Proof. Let $j : P \to Q$ be the canonical isomorphism, then $i_P = i_Q j$ by Remark 2.9. Now, since $Cof(\mathbf{C})$ [resp. $W(\mathbf{C})$] is closed under composition, Lemma 3.8 follows from Remark 3.4.

Proposition 3.9. Given a commutative cube in C of the form

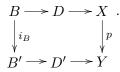


where the back face and the front face are pushout squares. Let P denote the pushout of the diagram $C \leftarrow A \rightarrow A'$ and let $i_P : P \rightarrow C'$ be the morphism induced by the left-hand face of (3.3). If i_B and i_P are cofibrations, then so is i_D .

Proof. By Proposition 3.6i), it's enough to find a lift in any given commutative diagram



where p is in $W \cap Fib$. Since i_B is a cofibration, there exists a lift $h_{B'} : B' \to X$ in



Defining $h_{A'}$ as the composition $A' \to B' \stackrel{h_{B'}}{\to} X$ yields a commutative diagram

$$A \xrightarrow{} A' \\ \downarrow \qquad \qquad \downarrow^{h_{A'}} \\ C \xrightarrow{} D \xrightarrow{} X$$

and hence an induced morphism $P \to X$. This morphism makes the diagram

$$\begin{array}{ccc} P & \longrightarrow & X \\ & & \downarrow_{i_P} & & \downarrow_{p} \\ C' & \longrightarrow & D' & \longrightarrow & Y \end{array} \tag{3.4}$$

commute, as one checks using the universal property of the pushout P. Since i_P is a cofibration by assumption, there exists a lift $h_{C'}: C' \to X$ in (3.4). It makes the square

$$\begin{array}{c} A' \longrightarrow B' \\ \downarrow & \downarrow \\ C' \xrightarrow{h_{C'}} X \end{array}$$

commute. This square induces a morphism h from the pushout D' to X. One checks that h is the desired lift, using the universal property of the pushouts D and D'.

3.2 The homotopy category of a model category

Definition 3.10. Let **C** be a category and W a class of morphisms of **C**. A *localization of* **C** with respect to W is a category **D** together with a functor $F : \mathbf{C} \to \mathbf{D}$ such that the following two conditions hold:

- i) F(f) is an isomorphism for every f in W.
- ii) If G is a functor from **C** to a category \mathbf{D}' , such that G(f) is an isomorphism for every f in W, then there exists a unique functor $G' : \mathbf{D} \to \mathbf{D}'$ with G'F = G.

Remark 3.11. Let **C** be a category and W a class of morphisms of **C**. If (\mathbf{D}, F) and (\mathbf{D}', F') are localizations of **C** with respect to W, then the unique functor G' such that G'F = F' is an isomorphism. Therefore, if a localization exists, we will sometimes speak of the localization. Let **C** be a model category until the end of this subsection. The localization (Ho(**C**), $\gamma_{\mathbf{C}}$) of **C** with respect to the class of weak equivalences W exists by Theorem 6.2 of [1]. This fact makes the following definition possible.

Definition 3.12. The homotopy category of the model category \mathbf{C} is the localization of \mathbf{C} with respect to W.

Until the end of this subsection let F be a functor from \mathbf{C} to another model category \mathbf{D} . The homotopy colimit functor will be defined as a total left derived functor, which is defined as follows.

Definition 3.13. A total left derived functor $\mathbf{L}F$ for the functor F is a functor $\mathbf{L}F : \operatorname{Ho}(\mathbf{C}) \to \operatorname{Ho}(\mathbf{D})$ together with a natural transformation $t : (\mathbf{L}F)\gamma_{\mathbf{C}} \to \gamma_{\mathbf{D}}F$ such that for any pair (G, s) of a functor $G : \operatorname{Ho}(\mathbf{C}) \to \operatorname{Ho}(\mathbf{D})$ and a natural transformation $s : G\gamma_{\mathbf{C}} \to \gamma_{\mathbf{D}}F$, there exists a unique natural transformation $s' : G \to \mathbf{L}F$ satisfying

$$t \circ s'|_{\gamma_{\mathbf{C}}} = s,$$

where the natural transformation $s'|_{\gamma_{\mathbf{C}}} : G\gamma_{\mathbf{C}} \to (\mathbf{L}F)\gamma_{\mathbf{C}}$ is given by $s'_{\gamma_{\mathbf{C}}(X)}$ in an object X of **C**.

Remark 3.14. Assume that $(\mathbf{L}F, t)$ and $(\mathbf{L}'F, t')$ are total left derived functors for F. Then the unique natural transformation $s' : \mathbf{L}'F \to \mathbf{L}F$ such that $t \circ s'|_{\gamma_{\mathbf{C}}} = t'$ is a natural equivalence. Therefore, if a total left derived functor for F exists, we will sometimes speak of the left derived functor.

The proof of the following theorem is given on page 114 in [1].

Theorem 1. Assume $G : \mathbf{D} \to \mathbf{C}$ is a functor, which is right adjoint to F and which carries morphisms of $Fib(\mathbf{D})$ to $Fib(\mathbf{C})$ and morphisms of $Fib(\mathbf{D}) \cap W(\mathbf{D})$ to $Fib(\mathbf{C}) \cap W(\mathbf{C})$. Then the total left derived functor $\mathbf{L}F : Ho(\mathbf{C}) \to Ho(\mathbf{D})$ for F exists.

4 Homotopy colimits

In this section let \mathbf{C} be a model category. We want to define the homotopy colimit functor as the total left derived functor for the functor colim : $\mathbf{C}^{\mathbf{D}} \rightarrow \mathbf{C}$. Therefore, we have to equip the functor category $\mathbf{C}^{\mathbf{D}}$ with a suitable model category structure, which can be done under the assumption that \mathbf{D} is as in the next definition.

Definition 4.1. A non-empty, finite category **D** is said to be *very small* if there exists an integer $N \ge 1$ such that for any composition $f_N \dots f_2 f_1$ of morphisms $(f_i)_{1 \le i \le N}$ in **D**, at least one morphism f_i is an identity morphism.

More geometrically, note that a non-empty, finite category **D** is very small if and only if it has no cycles, i.e. given any integer $n \ge 1$ and any composition $f_n...f_2f_1: d \to d$ of morphisms $(f_i)_{1 \le i \le n}$ in **D**, then each morphism f_i is the identity morphism id_d . The advantage of a very small category is that it enables us to do induction involving the degree, which is defined as follows.

Definition 4.2. Let **D** be a very small category and d any object of **D**. The *degree* deg(d) of d is defined by

 $\deg(d) := \max(\{0\} \cup \{n \ge 1; \text{ there exists a composition } f_n \dots f_2 f_1 : e \to d \text{ of}$ morphisms $(f_i \ne \mathrm{id}_d)_{1 \le i \le n} \text{ in } \mathbf{D}\}),$

the total degree $\deg(\mathbf{D})$ of \mathbf{D} by $\deg(\mathbf{D}) := \max(\{\deg(d); d \text{ an object of } \mathbf{D}\}).$

Remark 4.3. Let e, d be objects of a very small category **D** and assume, there exists a non-identity morphism $e \to d$. Then $\deg(e) < \deg(d)$.

4.1 A model category structure for C^D

Let \mathbf{D} be a very small category. The functor category $\mathbf{C}^{\mathbf{D}}$ can be given a model category structure in the following way.

Let d be an object of **D**. Recall that an object m of the over category $\mathbf{D} \downarrow d$ is given by a morphism $m : e \to d$ in **D** and that a morphism k in $\mathbf{D} \downarrow d$ from $m : e \to d$ to $m' : e' \to d$ is a morphism $k : e \to e'$ in **D** such that m'k = m. Denote by ∂d the full subcategory of $\mathbf{D} \downarrow d$ which contains all objects of $\mathbf{D} \downarrow d$ except id_d . Let the functor $j_d : \partial d \to \mathbf{D}$ be given by $(m : e \to d) \mapsto e$ on objects and $k \mapsto k$ on morphisms. Composing the induced functor $(\cdot)|_{\partial d} : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}^{\partial d}$ with colim : $\mathbf{C}^{\partial d} \to \mathbf{C}$ gives a functor

$$\partial_d := \operatorname{colim} \circ (\cdot) |_{\partial d} : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}.$$

Let X be an object of $\mathbf{C}^{\mathbf{D}}$. The natural transformation $s_d^X : X|_{\partial d} \to \Delta(X(d))$ given in an object $m : e \to d$ of ∂d by X(m), induces the morphism

$$\alpha^{-1}(s_d^X):\partial_d(X)\longrightarrow X(d),$$

where the bijection α comes from the adjunction (colim, $\Delta_{\partial d}, \alpha$). This induced morphism is natural. Indeed, given any morphism $f: X \to Y$ in $\mathbf{C}^{\mathbf{D}}$, we show that the diagram

$$\begin{array}{ccc} \partial_d(X) \longrightarrow X(d) & (4.1) \\ & & & \downarrow^{\partial_d(f)} & & \downarrow^{f_d} \\ \partial_d(Y) \longrightarrow Y(d) \end{array}$$

commutes or equivalently, that

$$\alpha(\alpha^{-1}(s_d^Y)\partial_d(f)) = \alpha(f_d\alpha^{-1}(s_d^X))$$
(4.2)

holds. Using the naturality of α , one concludes that the left-hand side of (4.2) equals $s_d^Y \circ f|_{\partial d}$, which is $Y(m) \circ f_e$ in an object $m : e \to d$, and that the right-hand side equals $\Delta(f_d) \circ s_d^X$, which is $f_d \circ X(m)$ in $m : e \to d$. Finally, the equation $Y(m) \circ f_e = f_d \circ X(m)$ holds, since f is a natural transformation by assumption. Define the functor δ_d as the composition

$$\delta_d : \mathbf{Mor}(\mathbf{C}^{\mathbf{D}}) \longrightarrow \mathbf{C}^{a \leftarrow b \rightarrow c} \stackrel{\text{colim}}{\longrightarrow} \mathbf{C},$$

where the first functor is given by

$$(f: X \to Y) \mapsto (\partial_d(Y) \xleftarrow{\partial_d(f)} \partial_d(X) \longrightarrow X(d)) \text{ on objects}, (s, s') \mapsto (\partial_d(s'), \partial_d(s), s_d) \text{ on morphisms}.$$

Given any morphism $f: X \to Y$ in $\mathbf{C}^{\mathbf{D}}$, consider that $\delta_d(f)$ is a pushout by definition and that therefore the commutative square (4.1) induces a morphism

$$i_d(f): \delta_d(f) \to Y(d),$$

which is natural. Indeed, given any morphism (s, s') in $Mor(\mathbb{C}^{\mathbf{D}})$ from $f: X \to X'$ to $g: Y \to Y'$, one checks that

$$\begin{array}{c|c} \delta_d(f) \xrightarrow{i_d(f)} X'(d) \\ \delta_d(s,s') & \downarrow s'_d \\ \delta_d(g) \xrightarrow{i_d(g)} Y'(d) \end{array}$$

commutes using the universal property of pushouts, the equation $s'_d f_d = g_d s_d$ which holds by assumption and the naturality of the morphism $\partial_d(X') \rightarrow X'(d)$. Using the above constructed morphisms $(i_d(f))_d$ for a morphism fin $\mathbf{C}^{\mathbf{D}}$, we give $\mathbf{C}^{\mathbf{D}}$ a model category structure.

Theorem 2. Define a morphism f of C^D to be in

- i) $W(\mathbf{C}^{\mathbf{D}})$, if f_d is in $W(\mathbf{C})$ for all objects d of \mathbf{D} ,
- ii) $Fib(\mathbf{C}^{\mathbf{D}})$, if f_d is in $Fib(\mathbf{C})$ for all objects d of \mathbf{D} , and to be in
- iii) $Cof(\mathbf{C}^{\mathbf{D}})$, if $i_d(f)$ is in $Cof(\mathbf{C})$ for all objects d of \mathbf{D} .

Then C^D together with these three classes is a model category.

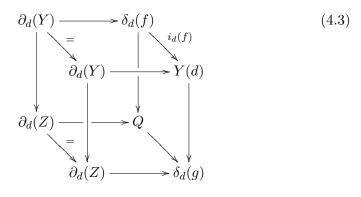
Remark 4.4. One can check directly, that the property of a morphism f in $\mathbf{C}^{\mathbf{D}}$ to be in $Cof(\mathbf{C}^{\mathbf{D}})$ doesn't depend on the choices of colimits involved in the construction of the morphisms $(i_d(f))_d$. However, this fact follows from Proposition 3.6i) after having proved Theorem 2.

Proof of Theorem 2. Since $W(\mathbf{C})$ and $Fib(\mathbf{C})$ are closed under composition and contain all identity morphisms, it follows immediately that $W(\mathbf{C}^{\mathbf{D}})$ and $Fib(\mathbf{C}^{\mathbf{D}})$ share the same properties.

To show that $Cof(\mathbf{C}^{\mathbf{D}})$ is closed under composition, let two morphisms $f: X \to Y, g: Y \to Z$ in $Cof(\mathbf{C}^{\mathbf{D}})$ be given. We have to prove, that $i_d(gf)$ is in $Cof(\mathbf{C})$ for all objects d of \mathbf{D} . By Lemma 3.8, it's enough to show that for any pushout Q of $\partial_d(Z) \leftarrow \partial_d(X) \to X(d)$, the morphism $i: Q \to Z(d)$ induced by the diagram

$$\begin{array}{c|c} \partial_d(X) \longrightarrow X(d) \\ \partial(g)\partial(f) = \partial(gf) \middle| & & & \downarrow (gf)_d = g_d f_d \\ \partial_d(Z) \longrightarrow Z(d) \end{array}$$

is a cofibration in **C**. Using Proposition 2.15, define such a Q as the pushout of $\partial_d(Z) \leftarrow \partial_d(Y) \rightarrow \delta_d(f)$. To show that the induced morphism $i: Q \rightarrow Z(d)$ is in $Cof(\mathbf{C})$, let $Q \rightarrow \delta_d(g)$ be the unique morphism such that



commutes. The pushout P of $\partial_d(Z) \leftarrow \partial_d(Y) \xrightarrow{=} \partial_d(Y)$ is just $\partial_d(Z)$ and the morphism $P \to \partial_d(Z)$ induced by the left-hand face of (4.3) is $\mathrm{id}_{\partial_d(Z)}$ and therefore in $Cof(\mathbf{C})$. Applying Proposition 3.9, we get that $Q \to \delta_d(g)$ is in $Cof(\mathbf{C})$. Finally, using the universal property of the pushout Q of $\partial_d(Z) \leftarrow \partial_d(X) \to X(d)$, one checks that i equals the composition

$$Q \to \delta_d(g) \stackrel{i_d(g)}{\to} Z(d)$$

and therefore is in $Cof(\mathbf{C})$ as a composition of two cofibrations in \mathbf{C} .

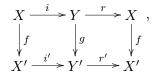
To prove that $Cof(\mathbf{C}^{\mathbf{D}})$ contains all identity morphisms of $\mathbf{C}^{\mathbf{D}}$, one has to show the following claim. For any object d of \mathbf{D} and any object X of $\mathbf{C}^{\mathbf{D}}$, the morphism $i_d(\mathrm{id}_X) : \delta_d(\mathrm{id}_X) \to X(d)$ is a cofibration in \mathbf{C} . Since $\partial_d(\mathrm{id}_X) = \mathrm{id}_{\partial_d(X)}$, it follows that X(d) is a pushout of $\partial_d(X) \leftarrow \partial_d(X) \to$ X(d). The claim is now a consequence of Lemma 3.8.

4.1.1 Proof of MC1-MC3

By MC1 in C and Proposition 2.13, it follows that any functor F from a finite category to $\mathbf{C}^{\mathbf{D}}$ has a limit. We want to show that F has a colimit or equivalently by Remark 2.14 that F^{op} has a limit. Note that $(\mathbf{C}^{\mathbf{D}})^{\text{op}}$ is isomorphic to $(\mathbf{C}^{\text{op}})^{\mathbf{D}^{\text{op}}}$ and conclude from Proposition 2.13 and MC1 in \mathbf{C}^{op} , that F^{op} has a limit. Hence, MC1 holds in $\mathbf{C}^{\mathbf{D}}$.

From MC2 for C, we will deduce MC2 for $\mathbf{C}^{\mathbf{D}}$. Let f, g be morphisms in $\mathbf{C}^{\mathbf{D}}$ such that gf is defined and such that two of the three morphisms f, g, gf are in $W(\mathbf{C}^{\mathbf{D}})$. Then for all objects d of \mathbf{D} , two of the three morphisms $f_d, g_d, g_d f_d$ are in $W(\mathbf{C})$. Thus by MC2 for C and the equality $g_d f_d = (gf)_d$, all three morphisms $f_d, g_d, (gf)_d$ are in $W(\mathbf{C})$ and hence f, g, gf are in $W(\mathbf{C}^{\mathbf{D}})$.

To show MC3, let $f: X \to X', g: Y \to Y'$ be morphisms in $\mathbb{C}^{\mathbf{D}}$ such that f is a retract of g, i.e. there exists a commutative diagram



such that $ri = id_X$ and $r'i' = id_{X'}$. Note that therefore f_d is a retract of g_d for all objects d of **D**. Thus the part of **MC3** dealing with fibrations [resp. weak equivalences] is a direct consequence of **MC3** for **C**. Assume that g is a cofibration, i.e. for all objects d of **D** the morphism $i_d(g)$ is in $Cof(\mathbf{C})$. We will deduce that f is a cofibration by showing that $i_d(f)$ is a retract of $i_d(g)$ and hence is in $Cof(\mathbf{C})$ by **MC3** for **C**. Consider the diagram

$$\delta_d(f) \xrightarrow{\delta_d(i,i')} \delta_d(g) \xrightarrow{\delta_d(r,r')} \delta_d(f)$$

$$\downarrow^{i_d(f)} \qquad \qquad \downarrow^{i_d(g)} \qquad \qquad \downarrow^{i_d(f)}$$

$$X'(d) \xrightarrow{i'_d} Y'(d) \xrightarrow{r'_d} X'(d)$$

The morphism $i_d(f)$ is natural as already shown, so the diagram commutes. By functoriality and since $ri = id_X$, $r'i' = id_{X'}$, it follows that the composition in the top row is the identity morphism. The equation $r'i' = id_{X'}$ implies that the composition in the bottom row equals $id_{X'(d)}$. Hence, $i_d(f)$ is a retract of $i_d(g)$.

4.1.2 Proof of MC4 and MC5

We will use induction over the total degree deg(D) of D to prove MC4, MC5 and the following proposition.

Proposition 4.5. Let $i : A \to B$ be a morphism in \mathbb{C}^{D} . Then i is in $Cof(\mathbb{C}^{D}) \cap W(\mathbb{C}^{D})$, if and only if $i_{d}(i)$ is in $Cof(\mathbb{C}) \cap W(\mathbb{C})$ for all objects d of D.

Remark 4.6. Note that to prove any direction of Proposition 4.5, it's enough to show that $\partial_d(i)$ is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ for all objects d of \mathbf{D} by Proposition 3.7ii) and **MC2** in \mathbf{C} .

For the proof of the initial case $\deg(\mathbf{D}) = 0$, we will use the following lemma.

Lemma 4.7. Assume deg(D) = 0. Let d be any object of D and $i : A \to B$ a morphism in C^{D} . Then $\partial_{d}(i)$ is an isomorphism in C. Furthermore, the morphism $i_{d}(i)$ is in Cof(C), if and only if i_{d} is in Cof(C).

Proof of Lemma 4.7. Since deg(**D**) = 0 implies deg(d) = 0, it follows by Remark 4.3 that ∂d is the empty category. Hence, the colimits $\partial_d(A)$ and $\partial_d(B)$ are initial objects and $\partial_d(i)$ is an isomorphism. A pushout of $\partial_d(B) \leftarrow$ $\partial_d(A) \to A(d)$ is given by A(d) and the morphism $A(d) \to B(d)$ induced by the commutative square

$$\begin{array}{c} \partial_d(A) \longrightarrow A(d) \\ & \swarrow \\ \partial_d(B) \longrightarrow B(d) \end{array}$$

is i_d . Hence, the second statement of Lemma 4.7 follows from Lemma 3.8.

We show Proposition 4.5, MC4 and MC5 in the initial case. By Remark 4.6, Lemma 4.7 implies Proposition 4.5, since any isomorphism in C is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ by Remark 3.4.

To prove MC4, let a commutative diagram

$$\begin{array}{ccc} A & \stackrel{f}{\longrightarrow} X \\ & & & \downarrow^{p} \\ \downarrow^{i} & & \downarrow^{p} \\ B & \stackrel{g}{\longrightarrow} Y \end{array}$$

in $\mathbf{C}^{\mathbf{D}}$ be given, such that *i* is in $Cof(\mathbf{C}^{\mathbf{D}})$ and *p* is in $Fib(\mathbf{C}^{\mathbf{D}})$. We have to find a lift $h: B \to X$, whenever *p* is in $W(\mathbf{C}^{\mathbf{D}})$ [resp. *i* is in $W(\mathbf{C}^{\mathbf{D}})$]. By Lemma 4.7, it follows that i_d is in $Cof(\mathbf{C})$ for all objects *d* of **D**. If *p* is in $W(\mathbf{C}^{\mathbf{D}})$ [resp. if *i* is in $W(\mathbf{C}^{\mathbf{D}})$], then applying **MC4**i) [resp. **MC4**ii)] in **C** yields an objectwise lift $h_d: B(d) \to X(d)$. These objectwise lifts $(h_d)_d$ fit together to give the desired lift $h: B \to X$.

To prove **MC5**i) [resp. **MC5**ii)], let $f : A \to B$ be a morphism in $\mathbf{C}^{\mathbf{D}}$. Using **MC5**i) [resp. **MC5**ii)] in **C**, factor for every object d of **D** the morphism f_d as $f_d = p_d i_d$, where i_d is in $Cof(\mathbf{C})$ [resp. $Cof(\mathbf{C}) \cap W(\mathbf{C})$] and p_d is in $W(\mathbf{C}) \cap Fib(\mathbf{C})$ [resp. $Fib(\mathbf{C})$]. By construction and Lemma 4.7, the morphisms $(i_d)_d$ fit together to give a morphism i in $Cof(\mathbf{C}^{\mathbf{D}})$

[resp. $Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$] and the morphisms $(p_d)_d$ define a morphism p in $W(\mathbf{C}^{\mathbf{D}}) \cap Fib(\mathbf{C}^{\mathbf{D}})$ [resp. $Fib(\mathbf{C}^{\mathbf{D}})$]. The factorization f = pi shows that **MC5**i) [resp. **MC5**ii)] holds.

To show the induction step, assume that $n := \deg(\mathbf{D}) \geq 1$. We prove Proposition 4.5. Let a morphism $i : A \to B$ in $\mathbf{C}^{\mathbf{D}}$ and any object d of \mathbf{D} be given. For the proof of any direction, it's enough to show that $\partial_d(i)$ is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ or equivalently by Proposition 3.6ii), to find a lift in any commutative diagram

$$\partial_d(A) \xrightarrow{f} C$$

$$\downarrow^{\partial_d(i)} \qquad \downarrow^p$$

$$\partial_d(B) \xrightarrow{g} D$$

$$(4.4)$$

in **C**, where p is in $Fib(\mathbf{C})$. From the commutativity of the diagram (4.4) and the definition of $\partial_d(i)$, it follows that also the diagram

$$A|_{\partial d} \longrightarrow \Delta(\partial_d(A)) \xrightarrow{\Delta(f)} \Delta(C)$$

$$\downarrow^{i|_{\partial d}} \qquad \qquad \downarrow^{\Delta(\partial_d(i))} \qquad \qquad \downarrow^{\Delta(p)}$$

$$B|_{\partial d} \longrightarrow \Delta(\partial_d(B)) \xrightarrow{\Delta(g)} \Delta(D)$$

$$(4.5)$$

in $\mathbf{C}^{\partial d}$ commutes. In any of the two directions of Proposition 4.5 we will apply the induction hypothesis to find a lift $h: B|_{\partial d} \to \Delta(C)$ in (4.5), which will induce the desired lift $\partial_d(B) \to C$ in (4.4). Indeed, assume h is a lift in (4.5) and let $h': \partial_d(B) \to C$ be the induced morphism, i.e. $h' = \alpha^{-1}(h)$, where the bijection α comes from the adjunction (colim, $\Delta_{\partial d}, \alpha$). Note that by the naturality of α the equations $\alpha(h'\partial_d(i)) = \alpha(h')i|_{\partial d}$ and $\alpha(ph') =$ $\Delta(p)\alpha(h')$ hold. Since $h = \alpha(h')$ is a lift, deduce that $\alpha(f) = \alpha(h'\partial_d(i))$ and $\alpha(g) = \alpha(ph')$. This shows that h' is a lift in (4.4). To find a lift h in (4.5), note that the category ∂d is very small with deg(∂d) < n and that $\Delta(p)$ is in $Fib(\mathbf{C}^{\partial d})$. Hence, using the induction hypothesis to apply **MC4**ii) in $\mathbf{C}^{\partial d}$, it's enough to show that $i|_{\partial d}$ is in $Cof(\mathbf{C}^{\partial d}) \cap W(\hat{\mathbf{C}}^{\partial d})$, that is by definition that $i_m(i|_{\partial d})$ is in $Cof(\mathbf{C})$ and that $(i|_{\partial d})_m$ is in $W(\mathbf{C})$ for every object $m: e \to d$ of ∂d . We calculate $i_m(i|_{\partial d})$. Recall that an object of the subcategory ∂m of $\partial d \downarrow m$ is given by a non-identity morphism $k: (m': e' \to d) \to m$ in ∂d , which by definition of ∂d is a non-identity morphism $k: e' \to e$ in **D** such that mk = m'. Furthermore, a morphism $g: k \to l$ in ∂m from $k: (m': e' \to d) \to m$ to $l: (m'': e'' \to d) \to m$ is a morphism $g: m' \to m''$ in ∂d such that lg = k, which is a morphism $g: e' \to e''$ in **D** with m''g = m' and $lg = k: e' \to e$. Hence, we can define a functor $j': \partial e \to \partial m$ by

$$(k:e' \to e) \mapsto (k:mk \to m)$$
 on objects and
 $(g:k \to l) \mapsto (g:j'(k) \to j'(l))$ on morphisms,

which is an isomorphism. Using (2.1), one concludes that the induced functor $(\cdot)|_{j'} : \mathbf{C}^{\partial m} \to \mathbf{C}^{\partial e}$ is an isomorphism. Since $\Delta_{\partial e} = (\cdot)|_{j'} \Delta_{\partial m}$ by Example 2.6, it follows by Remark 2.11 that the composition colim $\circ (\cdot)|_{j'}$ of colim : $\mathbf{C}^{\partial e} \to \mathbf{C}$ with $(\cdot)|_{j'}$ is the colimit functor $\mathbf{C}^{\partial m} \to \mathbf{C}$. Hence, the functor $\partial_m : \mathbf{C}^{\partial d} \to \mathbf{C}$ is the composition colim $\circ (\cdot)|_{j'} \circ (\cdot)|_{\partial m}$. The composition $j_d \circ j_m \circ j'$ of the functors $j' : \partial e \to \partial m$, $j_m : \partial m \to \partial d$, $j_d : \partial d \to \mathbf{D}$ equals j_e . It follows by (2.1) that $(\cdot)|_{j'} (\cdot)|_{\partial m} (\cdot)|_{\partial d} = (\cdot)|_{\partial e}$ and hence $\partial_m(A|_{\partial_d}) = \operatorname{colim}(A|_{\partial_e}) = \partial_e(A), \ \partial_m(B|_{\partial_d}) = \partial_e(B)$ and $\partial_m(i|_{\partial d}) = \partial_e(i)$. Note that $i|_{\partial d} : A|_{\partial d} \to B|_{\partial d}$ is just i_e in the object $m : e \to d$ of ∂d . It follows that the diagram

$$\begin{array}{c} \partial_m(A|_{\partial d}) \longrightarrow (A|_{\partial d})(m) \\ \downarrow^{\partial_m(i|_{\partial d})} & \downarrow^{(i|_{\partial d})m} \\ \partial_m(B|_{\partial d}) \longrightarrow (B|_{\partial d})(m) \end{array}$$

is just

$$\begin{array}{c} \partial_e(A) \longrightarrow A(e) \\ & \downarrow^{\partial_e(i)} & \downarrow^{i_e} \\ \partial_e(B) \longrightarrow B(e) \end{array}$$

by Remark 2.11, since for any Y in $\mathbf{C}^{\mathbf{D}}$, the natural transformation

$$s_m^{Y|_{\partial d}} : (Y|_{\partial d})|_{\partial m} \to \Delta_{\partial m}(Y|_{\partial d}(m))$$

is in an object $k : m' \to m$ of ∂m given by $Y|_{\partial d}(k) = Y(j_d(k)) = Y(k)$ and hence the equation $(s_m^{Y|_{\partial d}})|_{j'} = s_e^Y$ holds. It follows that $i_m(i|_{\partial d})$ equals $i_e(i)$. Assuming now first, that i is in $Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$, one concludes that for every object $m : e \to d$ of ∂d the morphism $i_m(i|_{\partial d})$ is in $Cof(\mathbf{C})$ and that $(i|_{\partial d})_m = i_e$ is in $W(\mathbf{C})$. This shows one direction of Proposition 4.5. For the other direction, assume that $i_{d'}(i)$ is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ for all objects d' of \mathbf{D} . It follows that for any object $m : e \to d$ in ∂d the morphism $i_m(i|_{\partial d})$ is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$. Using the induction hypothesis to apply Proposition 4.5 one deduces that $i|_{\partial d}$ is in $Cof(\mathbf{C}^{\partial d}) \cap W(\mathbf{C}^{\partial d})$, which completes the proof of Proposition 4.5.

In the proof of **MC4** and **MC5** we will use the following notation. Let \mathbf{D}^{n-1} be the full subcategory of **D** which contains precisely the objects e of **D** with deg $(e) \leq n-1$. It is very small and has total degree n-1. The inclusion functor $j: \mathbf{D}^{n-1} \to \mathbf{D}$ induces the functor $(\cdot)|_{\mathbf{D}^{n-1}}: \mathbf{C}^{\mathbf{D}} \to \mathbf{C}^{\mathbf{D}^{n-1}}$, which carries weak equivalences [resp. fibrations] in $\mathbf{C}^{\mathbf{D}}$ to weak equivalences [resp. fibrations] in $\mathbf{C}^{\mathbf{D}}$ to weak equivalences [resp. fibrations] in $\mathbf{C}^{\mathbf{D}(n-1)}$. Note that $\mathbf{D}^{n-1} \downarrow e = \mathbf{D} \downarrow e$ for any object e of **D** with deg(e) < n. One concludes that for any morphism f of $\mathbf{C}^{\mathbf{D}}$ the equation

$$i_e(f|_{\mathbf{D}^{n-1}}) = i_e(f)$$
 (4.6)

holds and in particular that $(\cdot)|_{\mathbf{D}^{n-1}}$ carries cofibrations in $\mathbf{C}^{\mathbf{D}}$ to cofibrations in $\mathbf{C}^{\mathbf{D}^{n-1}}$.

To show MC4i) [resp. MC4ii)], let a commutative diagram



in $\mathbf{C}^{\mathbf{D}}$ be given, such that i is in $Cof(\mathbf{C}^{\mathbf{D}})$ and p is in $Fib(\mathbf{C}^{\mathbf{D}})$. Assuming that p is in $W(\mathbf{C}^{\mathbf{D}})$ [resp. i is in $W(\mathbf{C}^{\mathbf{D}})$], one has to find a lift $h: B \to X$. Use the induction hypothesis to apply $\mathbf{MC4i}$ [resp. $\mathbf{MC4ii}$] to find a lift $h^{n-1}: B|_{\mathbf{D}^{n-1}} \to X|_{\mathbf{D}^{n-1}}$ in the commutative square

Now, the strategy is to find for each object d of degree n of \mathbf{D} an objectwise lift $B(d) \to X(d)$, such that these lifts and h^{n-1} fit together to give the desired lift $h: B \to X$. Let the natural transformation

$$B|_{\partial d} \to \Delta_{\partial d}(X(d))$$

be given by the composition $B(e) \xrightarrow{(h^{n-1})_e} X(e) \xrightarrow{X(m)} X(d)$ in an object $m: e \to d$ of ∂d . The induced morphism $\partial_d(B) \to X(d)$ makes the diagram

$$\begin{array}{c|c} \partial_d(A) \xrightarrow{\alpha^{-1}(s_d^A)} A(d) & (4.7) \\ \partial_{d}(i) & \downarrow f_d \\ \partial_d(B) \longrightarrow X(d) \end{array}$$

commute. Indeed, by the naturality of α , one concludes that $\alpha(f_d \alpha^{-1}(s_d^A)) = \Delta(f_d)s_d^A$, which is $f_d A(m)$ in an object $m : e \to d$ of ∂d , and that α of the composition $\partial_d(A) \to \partial_d(B) \to X(d)$ equals the composition $A|_{\partial_d} \to B|_{\partial_d} \to \Delta(X(d))$, which is $X(m)(h^{n-1})_e i_e$ in $m : e \to d$. Finally, the equation $f_d A(m) = X(m)(h^{n-1})_e i_e$ holds, since $h^{n-1}i|_{\mathbf{D}^{n-1}} = f|_{\mathbf{D}^{n-1}}$ and since f is a natural transformation by assumption. Let $\delta_d(i) \to X(d)$ be the morphism induced by the commutative diagram (4.7). Using the universal property of pushouts, one checks that it makes

$$\begin{array}{cccc}
\delta_d(i) &\longrightarrow X(d) & (4.8) \\
\downarrow^{i_d(i)} & & \downarrow^{p_d} \\
B(d) &\xrightarrow{g_d} & Y(d)
\end{array}$$

commute. Apply **MC4**i) in **C** [resp. apply **MC4**ii) in **C** and Proposition 4.5] to find a lift $h_d : B(d) \to X(d)$ in (4.8). One checks that any lift $B(d) \to X(d)$ in (4.8) is also a lift in

$$\begin{array}{c|c} A(d) & \xrightarrow{f_d} & X(d) \\ \vdots \\ i_d & \downarrow \\ B(d) & \xrightarrow{g_d} & Y(d) \end{array}$$

The desired lift $h: B \to X$ can now be defined as $(h^{n-1})_e$ in an object e of **D** with deg(e) < n and as the constructed lift h_d of (4.8) in an object d of degree n of **D**. To show that h is a natural transformation, note that by Remark 4.3 and since h^{n-1} is a natural transformation, one only has to consider morphisms $m: e \to d$ in **D**, where deg(e) < n, deg(d) = n, and to check that $X(m)(h^{n-1})_e = h_d B(m)$ holds. Denoting colim $(B|_{\partial_d}) = (\partial d(B), t)$, this is done by the following sequence of equations of compositions of morphisms:

$$B(e) \xrightarrow{B(m)} B(d) \xrightarrow{h_d} X(d) = B(e) \xrightarrow{t_m} \partial_d(B) \longrightarrow B(d) \xrightarrow{h_d} X(d)$$
$$= B(e) \xrightarrow{t_m} \partial_d(B) \longrightarrow \delta_d(i) \xrightarrow{i_d(i)} B(d) \xrightarrow{h_d} X(d)$$
$$= B(e) \xrightarrow{t_m} \partial_d(B) \longrightarrow \delta_d(i) \longrightarrow X(d)$$
$$= B(e) \xrightarrow{t_m} \partial_d(B) \longrightarrow X(d)$$
$$= B(e) \xrightarrow{(h^{n-1})_e} X(e) \xrightarrow{X(m)} X(d).$$

To prove **MC5**, we have to factor a given morphism $f : A \to B$ in $\mathbb{C}^{\mathbf{D}}$ in the two ways i) and ii). Use the induction hypothesis to factor $f|_{\mathbf{D}^{n-1}}$ as

$$A|_{\mathbf{D}^{n-1}} \xrightarrow{h^{n-1}} X^{n-1} \xrightarrow{g^{n-1}} B|_{\mathbf{D}^{n-1}}$$

$$(4.9)$$

as in **MC5**i) [resp. **MC5**ii)]. Let d be an object of degree n of **D**. Note that the functor $j_d : \partial d \to \mathbf{D}$ carries any object m of ∂d to an object $j_d(m)$ with $\deg(j_d(m)) < n$. It therefore induces a functor $j'_d : \partial d \to \mathbf{D}^{n-1}$, which composed with the inclusion functor $j : \mathbf{D}^{n-1} \to \mathbf{D}$ equals j_d . For the induced functors follows

$$(\cdot)|_{\partial d} = (\cdot)|_{j'_d} \circ (\cdot)|_{\mathbf{D}^{n-1}}$$

$$(4.10)$$

by (2.1). Hence, applying the composition $\operatorname{colim} \circ (\cdot)|_{j'_d}$ of $\operatorname{colim} : \mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ with $(\cdot)|_{j'_d}$ to (4.9) yields a factorization

$$\partial_d(A) \to \operatorname{colim}(X^{n-1}|_{j'_d}) \to \partial_d(B)$$
 (4.11)

of $\partial_d(f)$. Now define X(d) and h_d through the pushout square

$$\begin{array}{ccc} \partial_d(A) & \longrightarrow A(d) \ . & (4.12) \\ & & & & \downarrow^{h_d} \\ \operatorname{colim}(X^{n-1}|_{j'_d}) & \longrightarrow X(d) \end{array}$$

Denote $\operatorname{colim}(X^{n-1}|_{j'_d}) = (\operatorname{colim}(X^{n-1}|_{j'_d}), t^d)$ and for a non-identity morphism $m : e \to d$ define the morphism $X(m) : X^{n-1}(e) \to X(d)$ as the composition

$$X^{n-1}(e) \xrightarrow{(t^d)_m} \operatorname{colim}(X^{n-1}|_{j'_d}) \to X(d).$$
(4.13)

Furthermore, set $X(\operatorname{id}_d) = \operatorname{id}_{X(d)}$. Then these choices, X^{n-1} and h^{n-1} fit together to give a functor $X : \mathbf{D} \to \mathbf{C}$ and a natural transformation $h : A \to X$. Indeed, let $m : e \to d$ be a non-identity morphism in \mathbf{D} with $\operatorname{deg}(d) = n$. If $m' : e' \to e$ is another non-identity morphism in \mathbf{D} , then one deduces by the naturality of t^d that $(t^d)_m X^{n-1}(m') = (t^d)_{mm'}$ holds and therefore X(mm') = X(m)X(m'). To prove the naturality of h, check that

$$\begin{array}{c} A(e) \xrightarrow[(h^{n-1})_e]{} X(e) \\ A(m) \downarrow \qquad \qquad \downarrow X(m) \\ A(d) \xrightarrow{h_d} X(d) \end{array}$$

commutes by noting that $(h^{n-1})_e = (h^{n-1}|_{j'_d})_m$ and by using the commutativity of the diagram (4.12) and of the diagram

We want to show that h is in $Cof(\mathbf{C}^{\mathbf{D}})$ [resp. h is in $Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$]. Note that $X|_{\mathbf{D}^{n-1}} = X^{n-1}$ and $h|_{\mathbf{D}^{n-1}} = h^{n-1}$ by definition. In case i), since by construction h^{n-1} is in $Cof(\mathbf{C}^{\mathbf{D}^{n-1}})$, it follows by (4.6) that $i_e(h)$ is in $Cof(\mathbf{C})$ for every object e of \mathbf{D} with deg(e) < n. Similarly, in case ii), since by construction h^{n-1} is in $Cof(\mathbf{C}^{\mathbf{D}^{n-1}}) \cap W(\mathbf{C}^{\mathbf{D}^{n-1}})$, it follows by Proposition 4.5 and by (4.6) that $i_e(h)$ is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ for every object e of \mathbf{D} with deg(e) < n. Let d be an object of degree n of \mathbf{D} . By (4.10), deduce the equations $X|_{\partial d} = X^{n-1}|_{j'_d}$, $h|_{\partial d} = h^{n-1}|_{j'_d}$ and therefore $\partial_d(X) = \operatorname{colim}(X^{n-1}|_{j'_d}), \partial_d(h) = \operatorname{colim}(h^{n-1}|_{j'_d})$. Recall that the morphism $\partial_d(X) \to X(d)$ is induced by the natural transformation s_d^X which is given in an object $m : e \to d$ of ∂d by X(m). Since X(m) is the composition (4.13), it follows that $\partial_d(X) \to X(d)$ is just $\operatorname{colim}(X^{n-1}|_{j'_d}) \to X(d)$. Hence, $\delta_d(h) = X(d)$ and $i_d(h)$ equals $\operatorname{id}_{X(d)}$ and in particular is in $\operatorname{Cof}(\mathbf{C}) \cap W(\mathbf{C})$. Thus in case i) we have shown that h is in $\operatorname{Cof}(\mathbf{C}^{\mathbf{D}})$ and by Proposition 4.5, it follows in case ii) that h is in $\operatorname{Cof}(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$.

Using g^{n-1} , we want to define a natural transformation $g: X \to B$ such that f = gh. Let d be an object of degree n of \mathbf{D} . The natural transformation $X|_{\partial d} \to \Delta(B(d))$ given by $B(m)(g^{n-1})_e$ in an object $m: e \to d$ of ∂d , induces a morphism $\partial_d(X) \to B(d)$ in \mathbf{C} . One checks that this morphism makes the square

$$\begin{array}{c|c} \partial_d(A) \longrightarrow A(d) \\ \partial_d(h) & & f_d \\ \partial_d(X) \longrightarrow B(d) \end{array}$$

commute. This gives an induced morphism g_d from the pushout X(d) to B(d) such that $f_d = g_d h_d$. The constructed morphisms $(g_d)_d$ and g^{n-1} fit together to give a natural transformation $g: X \to B$. Indeed, if $m: e \to d$ is a non-identity morphism in **D** with $\deg(d) = n$, then $B(m)g_e = g_d X(m)$ holds, as is shown by the following sequence of equations of compositions of morphisms:

$$X(e) \stackrel{(g^{n-1})_e}{\to} B(e) \stackrel{B(m)}{\to} = X(e) \stackrel{(t^d)_m}{\to} \partial_d(X) \to B(d)$$
$$= X(e) \stackrel{(t^d)_m}{\to} \partial_d(X) \to X(d) \stackrel{g_d}{\to} B(d)$$
$$= X(e) \stackrel{X(m)}{\to} X(d) \stackrel{g_d}{\to} B(d).$$

We have factored f as gh, such that h is in $Cof(\mathbf{C}^{\mathbf{D}})$ [resp. $Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$]. Thus to prove **MC5**i) and **MC5**ii), it's enough to factor g as pi in $\mathbf{C}^{\mathbf{D}}$ such that respectively, $i \in Cof(\mathbf{C}^{\mathbf{D}})$, $p \in W(\mathbf{C}^{\mathbf{D}}) \cap Fib(\mathbf{C}^{\mathbf{D}})$ and $i \in Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$, $p \in Fib(\mathbf{C}^{\mathbf{D}})$. Set $Z^{n-1} := X|_{\mathbf{D}^{n-1}} : \mathbf{D}^{n-1} \to \mathbf{C}$, let $i^{n-1} : X|_{\mathbf{D}^{n-1}} \to Z^{n-1}$ be the identity morphism and define the natural transformation $p^{n-1} := g^{n-1} : Z^{n-1} \to B|_{\mathbf{D}^{n-1}}$. Let d be any object of degree n of \mathbf{D} . Using **MC5** in \mathbf{C} , factor g_d as

$$X(d) \xrightarrow{i_d} Z(d) \xrightarrow{p_d} B(d),$$

as in **MC5**i) [resp. **MC5**ii)]. For any non-identity morphism $m : e \to d$ in **D** set $Z(m) := i_d X(m) : Z(e) \to Z(d)$ and let the morphism $Z(\mathrm{id}_d)$ in **C** be given by $\mathrm{id}_{Z(d)}$. Then these choices and Z^{n-1} fit together to give a functor $Z : \mathbf{D} \to \mathbf{C}$. Indeed, if $m' : e' \to e, m : e \to d$ are non-identity morphisms in **D** such that $\mathrm{deg}(d) = n$, then

$$Z(mm') = i_d X(mm') = i_d X(m) X(m') = Z(m) Z(m').$$

The constructed morphisms $(i_d)_d$ and i^{n-1} fit together to give a natural transformation $i: X \to Z$. Indeed, given a non-identity morphism $m: e \to d$ in **D** with deg(d) = n, one has $i_e = \operatorname{id}_{Z(e)}$ and thus $Z(m)i_e = i_d X(m)$. The morphisms $(p_d)_d$ and p^{n-1} yield a natural transformation $p: Z \to B$. Indeed, for a non-identity morphism $m: e \to d$ in **D** with deg(d) = n holds

$$B(m)p_e = B(m)g_e = g_d X(m) = p_d i_d X(m) = p_d Z(m).$$

Note that for an object d of arbitrary degree of \mathbf{D} , the functor $X|_{\partial d}$ equals $Z|_{\partial d}$ and that $i|_{\partial d} = \operatorname{id}_{X|_{\partial d}}$. Hence a pushout of $\partial_d(Z) \stackrel{=}{\leftarrow} \partial_d(X) \to X(d)$ is given by X(d) and the morphism $X(d) \to Z(d)$ induced by the commutative square

is just i_d . By Lemma 3.8, it follows that $i_d(i)$ is a cofibration [resp. weak equivalence] if and only if i_d is a cofibration [resp. weak equivalence]. Furthermore, note that if $\deg(d) < n$, then $i_d = \operatorname{id}_{Z(d)}$ and hence is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$.

In case i), since by construction for any object d of degree n of \mathbf{D} the morphism i_d is in $Cof(\mathbf{C})$, it follows that i is in $Cof(\mathbf{C}^{\mathbf{D}})$. One checks that by construction p is in $Fib(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$. Similarly, in case ii), the morphism i_d is in $Cof(\mathbf{C}) \cap W(\mathbf{C})$ for every object d of degree n of \mathbf{D} . By Proposition 4.5, it follows that $i \in Cof(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$. One checks that $p \in Fib(\mathbf{C}^{\mathbf{D}})$. This completes the proof of $\mathbf{MC5}$ and hence the induction step.

We have shown Theorem 2.

4.2 The homotopy colimit functor

Let **D** be a very small category. Let $\mathbf{C}^{\mathbf{D}}$ be the model category of Theorem 2. Recall that by Remark 2.10 the functor $\Delta : \mathbf{C} \to \mathbf{C}^{\mathbf{D}}$ is right adjoint to colim : $\mathbf{C}^{\mathbf{D}} \to \mathbf{C}$ and note that it carries morphisms of $Fib(\mathbf{C})$ and morphisms of $Fib(\mathbf{C}) \cap W(\mathbf{C})$ to $Fib(\mathbf{C}^{\mathbf{D}})$ and $Fib(\mathbf{C}^{\mathbf{D}}) \cap W(\mathbf{C}^{\mathbf{D}})$ respectively. By Theorem 1 we are finally in position to define homotopy colimits or more precisely, the homotopy colimit functor.

Definition 4.8. The homotopy colimit functor is the total left derived functor Lcolim : $Ho(\mathbf{C}^{\mathbf{D}}) \to Ho(\mathbf{C})$ for the functor colim : $\mathbf{C}^{\mathbf{D}} \to \mathbf{C}$.

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