

The drift criterion for recurrence

Let P be an irreducible Markov kernel on (X, E)

Recall that a function $V: X \rightarrow [0, \infty)$ is said to be unbounded off small sets if the small sets

$$\{x \mid V(x) \leq c\} \quad c \in (0, \infty)$$

all are small sets.

Theorem Let P be irreducible. Suppose that there exists a function $V: X \rightarrow [0, \infty)$ which is measurable and unbounded off small sets, and which satisfies that

$$\Delta V(x) \leq 0 \quad \forall x \notin C$$

for some small set $C \in \mathcal{E}^+$. Then P is recurrent.

Proof: Assume that the conclusion is wrong, so that P is transient.

As

$$L_x(C) = P_x(C) + \int_C L_y(C) dP_x(y)$$

We conclude that if

$$(*) \quad L_y(C) = 1 \quad \forall y \in C$$

$$\text{then} \quad L_x(C) = 1 \quad \forall x \in C$$

This implies that C is a recurrent \mathbb{E}^+ -set which is also small. But transience implies that any small \mathbb{E}^+ -set must be uniformly transient.

Hence we conclude that $(*)$ must be false. In other words: we can find some $x^* \notin C$ such that $L_{x^*}(C) < 1$.

$$\text{Pick } M > \frac{V(x^*)}{1 - L_{x^*}(C)}$$

observe that

$$C_M = \{x \mid \forall k \leq M\}$$

is small. If necessary we can choose an even larger M , and obtain that $C_M \in \mathbb{E}^+$. Combined with the assumption of transience, we have that C_M is uniformly transient.

Define a new Markov kernel $Q = (Q_x)_{x \in X}$ by

$$Q_x = \begin{cases} P_x & x \notin C \\ E_x & x \in C \end{cases}$$

If X_0, X_1, \dots is a P -Markov chain and τ is the first hitting time of C , then

$$X_{\tau+0}, X_{\tau+1}, X_{\tau+2}, \dots$$

is a time-homogeneous Markov chain with one-step transition probability Q .

• We note three properties of the Q -kernel

① If $x \notin C$, then

$$\int v(y) dQ_x(y) = \int v(y) dP_x(y) \leq v(x)$$

If $x \in C$, then

$$\int v(y) dQ_x(y) = \int v(y) dE_x(y) = v(x)$$

so actually

$$\int v(y) dQ_x(y) \leq v(x) \quad \forall x \in X$$

when the P -kernel only had this property outside C .

• ② If $x \notin C$ then $Q_x^n(C) \rightarrow L_x(C)$ for $n \rightarrow \infty$
The easiest proof of this is probably probabilistic:
let X_0, X_1, \dots be a P -chain with $X_0 \equiv x$. Let
 T be the first return time to C . Then

$$Q_x^n(C) = P(X_{0:n} \in C) = P(T \leq n) \rightarrow P(T < \infty)$$

and this last probability equals $L_x(C)$ per definition.

③ If $x \notin C$ and $A \subseteq C^c$ we have that

$$Q_x^y(A) \leq P_x^y(A)$$

We can again argue probabilistic. Let X_0, X_1, X_2, \dots be a P -chain with $X_0 = x$ and let τ be the first return time to C . Then

$$\begin{aligned} Q_x^y(A) &= P(X_{\tau+n} \in A) = P(X_n \in A, \tau > n) \\ &\leq P(X_n \in A) = P_x^y(A) \end{aligned}$$

To finish off the proof and arrive at a contradiction, we iterate on the claim in ① to get

$$\begin{aligned} V(x) &\geq \int V(y) dQ_x^y(y) \\ &\geq \int_{C^c \cap C_M^c} V(y) dQ_x^y(y) \\ &\geq M Q_x^y(C^c \cap C_M^c) \\ &= M (1 - Q_x^y(C \cup C_M)) \\ &= M (1 - Q_x^y(C) - Q_x^y(C_M \setminus C)) \end{aligned}$$

Observe that for $x \notin C$

$$Q_x^n(C_m \setminus C) \leq P_x^n(C_m \setminus C) \leq P_x^n(C_m)$$

As C_m is uniformly transient, there is $K < \infty$

$$\sum_{n=1}^{\infty} P_x^n(C_m) \leq K \quad \forall x \in \mathbb{T}$$

In particular $P_x^n(C_m) \rightarrow 0$.

So for $x \notin C$ we see that

$$Q_x^n(C) \rightarrow L_x(C), \quad Q_x^n(C_m \setminus C) \rightarrow 0$$

Hence we conclude that

$$V(x) \geq M(1 - L_x(C)) \quad \forall x \notin C$$

But M is specifically chosen such that this inequality is false for x^* .

This contradiction can only be resolved by refuting the original assumption that P is transient.

□