## Lecture 9: Axiomatic Set Theory

December 16, 2014

## Today

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- The iterative concept of set.
- The language of set theory (LOST).
- The axioms of Zermelo-Fraenkel set theory (ZFC).
- Justification of the axioms based on the iterative concept of set.


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for the negation of $x=y$ and $x \in y$, respectively.

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We will allow ourselves to use the standard abbreviations: $\vee, \wedge, \leftrightarrow$ and $\exists$ for "or", "and", "if and only if", and "there exists".

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(Sets that have the same members are identical.)

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(Every non-empty set has a member which has no members in common with it.)

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\forall w_{1} \cdots \forall w_{n} \forall z & (\forall x(x \in z \rightarrow \exists!y \varphi) \\
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(For any set $z$ and relation $R$ (which must be expressible by some first order formula $\varphi$ of LOST), if each member $x$ of $z$ bears the relation $R$ to exactly one set $y_{x}$, then there is a set to which all these $y_{x}$ belong.)

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(There is a set that has the empty set as a member and is closed under the operation $\mathcal{S}$.)

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For the next axiom, let " $v_{3} \subseteq v_{1}$ " abbreviate $" \forall v_{4}\left(v_{4} \in v_{3} \rightarrow v_{4} \in v_{1}\right)$ ".

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(For any set, there is a set to which all subsets of that set belong.)

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(If $x$ is a set of pairwise disjoint non-empty sets, then there is a set that has exactly one member in common with each member of $x$.)

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GROUP 2: Axiom of Union:

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GROUP 3: Axiom of Infinity:

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