

Classification of C^* -algebras associated to minimal dynamics

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Ss

Definition A *shift space* is a subset of $\mathfrak{a}^{\mathbb{Z}}$ which is closed in the product topology and under the shift map

$$\sigma((x_n)_{n\in\mathbb{Z}}) = (x_{n+1})_{n\in\mathbb{Z}}.$$

The alphabet a is a finite set.

A shift space is *of finite type* when it can be specified by a finite list of forbidden words.

Example
$$\underline{X}_{\{11,21\}} \ni \cdots 00100122200010201012 \cdots$$

Ss

Example With primitive substitutions such as

$$\tau(0) = 01$$
 $\tau(1) = 0$

one gets a shift space

$$\underline{\mathsf{X}}_{\tau} = \overline{\{\sigma^m(u) \mid m \in \mathbb{Z}\}}$$

where u is any periodic point of τ , i.e. $\tau^n(u) = u$.

In this case we could take n=2 and

 $u = \cdots 01001001.01001010010010100101001\cdots$

Some substitutions

$$\tau_{1}(\aleph) = \aleph \exists \aleph \qquad \tau_{1}(\exists) = \exists \aleph \aleph \exists$$

$$\tau_{2}(\alpha) = \alpha\beta \qquad \tau_{2}(\beta) = \alpha\beta\gamma\delta\epsilon \qquad \tau_{2}(\gamma) = \alpha\beta$$

$$\tau_{2}(\delta) = \gamma\delta\epsilon \qquad \tau_{2}(\epsilon) = \alpha\beta\gamma\delta\epsilon$$

$$\tau_{3}(1) = 1212345$$

$$\tau_{3}(2) = 12123451234512345$$

$$\tau_{3}(3) = 1212345 \qquad \tau_{3}(4) = 1234512345$$

$$\tau_{3}(5) = 12123451234512345$$

$$\tau_{4}(a) = ababacb \qquad \tau_{4}(b) = ababacbabacbabacb$$

$$\tau_{4}(c) = abacbabacb$$



Let X^+ denote the projection of \underline{X} down on $\mathfrak{a}^{\mathbb{N}}$. We set for finite words u,v written with letters from \mathfrak{a}

$$C(u|v) = \{ux \in X^+ \mid vx \in X^+\}$$

and define a C^* -algebra \mathcal{O}_X by generators $(S_a)_{a \in \mathfrak{a}}$ and relations

$$a \in \mathfrak{a} : S_a S_a^* S_a = S_a$$

$$C(u|v) = \bigcap_{i=1}^n C(u_i|v_i) : S_u S_v^* S_v S_u^* = \prod_{i=1}^n S_{u_i} S_{v_i}^* S_{v_i} S_{u_i}^*$$

$$C(u|v) = \bigsqcup_{i=1}^n C(u_i|v_i) : S_u S_v^* S_v S_u^* = \sum_{i=1}^n S_{u_i} S_{v_i}^* S_{v_i} S_{u_i}^*$$

with "∐" denoting disjoint unions. [Matsumoto]

Ss
$$\longrightarrow$$
 C^*

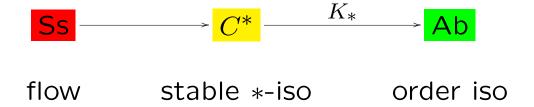
Key features:

- Matsumoto algebras associated to shifts of finite type are Cuntz-Krieger algebras.
- [Matsumoto, Carlsen]

$$\underline{\mathsf{X}} \simeq_{flow} \underline{\mathsf{Y}} \Longrightarrow \mathcal{O}_{\underline{\mathsf{X}}} \otimes \mathbb{K} \simeq \mathcal{O}_{\underline{\mathsf{Y}}} \otimes \mathbb{K}$$

Complication:

• Not always simple, even for minimal shift spaces like X_{τ} .



When are these maps injective?

This question has been successfully worked out in the case of irreducible shifts of finite type. [Cuntz, Krieger, Bowen, Franks, Rørdam]

The case for

$$\mathcal{O}_{\tau} := \mathcal{O}_{\mathsf{X}_{\tau}}$$

is stucturally very different but shares the property of the shift space being finitely presented. Work by Carlsen-E shows that not both maps are injective.

Good and bad news

- :-(\mathcal{O}_{τ} is never simple.
- [Carlsen] $0 \longrightarrow \mathbb{K}^{\mathsf{n}_{\mathcal{T}}} \longrightarrow \mathcal{O}_{\mathcal{T}} \longrightarrow C(\underline{\mathsf{X}}_{\mathcal{T}}) \rtimes_{\sigma} \mathbb{Z} \longrightarrow 0$
- [Carlsen-E] Computable K-theory:

$$K_{1}(\mathcal{O}_{\tau}) \longrightarrow K_{1}(C(\underline{X}_{\tau}) \rtimes_{\sigma} \mathbb{Z}) \longrightarrow K_{0}(\mathbb{K}^{\mathsf{n}_{\tau}}) \longrightarrow K_{0}(\mathcal{O}_{\tau}) \longrightarrow K_{0}(C(\underline{X}_{\tau}) \rtimes_{\sigma} \mathbb{Z})$$

$$0 \longrightarrow \mathbb{Z} \longrightarrow \mathbb{Z}^{n_{\tau}} \longrightarrow \mathsf{DG}(\tilde{\mathbf{A}}_{\tau}) \longrightarrow \mathsf{DG}(\mathbf{A}_{\tau})$$

- [Putnam] $C(\underline{X}_{\tau}) \rtimes_{\sigma} \mathbb{Z}$ is simple, AT and of real rank zero. \mathcal{O}_{τ} has real rank zero.
- \mathcal{O}_{τ} is not stably finite, nor purely infinite.

Bootstrap classification

Theorem [Rørdam 1994]

Let A and B be C^* -algebras each having one essential ideal I and J such that I,J,A/I and B/J are stable Kirchberg algebras satisfying the UCT. Then $A \simeq B$ precisely when

Unfortunately, this has proven very difficult to generalize to more ideals. Restorff is able to do the next case $0 \triangleleft I_0 \triangleleft I_1 \triangleleft A$, but by completely different methods.

Another bootstrap classification

Theorem [E-Restorff-Ruiz]

Let E_1 and E_2 be C^* -algebras each having an ideal B_i such that B_i are stable AF algebras and $A_i = E_i/B_i$ are simple AT algebras of real rank zero. Then $E_1 \otimes \mathbb{K} \simeq E_2 \otimes \mathbb{K}$ precisely when

$$K_1(E_1) \longrightarrow K_1(A_1) \longrightarrow K_0(B_1) \longrightarrow K_0(E_1) \longrightarrow K_0(A_1)$$

 $\simeq \mid \qquad \simeq \mid \qquad \simeq \mid \qquad \simeq \mid \qquad \simeq \mid$
 $K_1(E_2) \longrightarrow K_1(A_2) \longrightarrow K_0(B_2) \longrightarrow K_0(E_2) \longrightarrow K_0(A_2)$



Combine with results by Carlsen-E to get

Corollary $\mathcal{O}_{\tau} \otimes \mathbb{K} \simeq \mathcal{O}_{v} \otimes \mathbb{K}$ precisely when

$$\mathbb{Z} \longrightarrow \mathbb{Z}^{n_{\tau}} \longrightarrow K_0(\mathcal{O}_{\tau}) \longrightarrow K_0(C(\underline{\mathsf{X}}_{\tau}) \rtimes_{\sigma} \mathbb{Z})$$
 $\parallel \simeq \mid \simeq \mid \simeq \mid \simeq \mid$
 $\mathbb{Z} \longrightarrow \mathbb{Z}^{n_{v}} \longrightarrow K_0(\mathcal{O}_{v}) \longrightarrow K_0(C(\underline{\mathsf{X}}_{v}) \rtimes_{\sigma} \mathbb{Z})$

Corollary
$$\mathcal{O}_{\underline{\mathsf{X}}_{\tau}} \otimes \mathbb{K} \simeq \mathcal{O}_{\underline{\mathsf{X}}_{\upsilon}} \otimes \mathbb{K} \not\Longrightarrow \underline{\mathsf{X}}_{\tau} \simeq_{flow} \underline{\mathsf{X}}_{\upsilon}$$

Corollary
$$K_0(\mathcal{O}_{\underline{\mathsf{X}}_{\tau}}) \simeq K_0(\mathcal{O}_{\underline{\mathsf{X}}_{\upsilon}}) \not\Longrightarrow \mathcal{O}_{\tau} \otimes \mathbb{K} \simeq \mathcal{O}_{\upsilon} \otimes \mathbb{K}$$

Sketch of proof

Assume

① WLOG we may assume $A_1 = A_2 = A$, $B_1 = B_2 = B$, $\alpha_* = \mathrm{id}$, $\beta_* = \mathrm{id}$. [Strong classification/Elliott²]

Sketch of proof, continued

② Then with $x_i \in KK^1(A,B)$ representing the extensions we have

$$x_1b = ax_2$$

for some $a \in KK(A,A)^{-1}$, $b \in KK(B,B)^{-1}$. [Generalization of method by Rørdam]

- ③ $b = KK(id_B)$. And WLOG we may assume that also $a = KK(id_A)$ [Kishimoto-Kumjian]
- 4 Now $x_1=x_2$. Thus $au_1\oplus au_0\sim au_2\oplus au_0$. WLOG we may assume $au_1\sim au_2$ [Kucerovsky-Ng]
- ⑤ Done.

Good and bad news

- :- (A weak classification result
- :- (Range of the invariant currently unknown.
- :-) Similar methods apply to

$$0 \longrightarrow B \otimes \mathbb{K} \longrightarrow E \longrightarrow A \longrightarrow 0$$

if A and B both are simple unital AT algebras of real rank zero.

:-) Similar methods apply to

$$0 \longrightarrow B \longrightarrow E \longrightarrow A \longrightarrow 0$$

if A and B both are simple unital nuclear algebras of tracial rank zero with finitely generated K-theory [Dadarlat].